Relevant Context Inference
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Abstract

Relevant context inference (**RCI**) is a modular technique for flow- and context-sensitive data-flow analysis of statically typed object-oriented programming languages such as C^{++} and Java. *RCI* can be used to analyze complete programs as well as incomplete programs such as libraries; this approach does not require that the entire program be memory-resident during the analysis. We show that *RCI* can handle exceptions. We also discuss application of *RCI* to unit testing of libraries and explain how the information computed by *RCI* can be used for generating relevant test cases. *RCI* is presented in the context of *points-to analysis*. The empirical evidence obtained from a prototype implementation argues the effectiveness of *RCI*.

1 Introduction

Points-to analysis [EGH94] for statically typed object-oriented programming languages (e.g., Java, C^{++}) determines, at each program point, the objects to which a pointer may point during execution. This information is crucial to many applications, including static resolution of dynamically dispatched calls, side-effect analysis, data-flow-based testing, program slicing and aggressive compiler optimizations. The solution of *concrete type inference* [PC94], necessary for object-oriented optimizations such as method specialization and inlining, is subsumed by the solution of points-to analysis, because the concrete type of a pointer is the set of classes corresponding to the objects possibly pointed to by that pointer.

The goal of our analysis is to preserve precision as much as possible, without sacrificing scalability. Flow and context sensitivity affect both the precision and cost of analyses. A *flow-insensitive* algorithm ignores the ordering of statements within a method; by contrast, a *flow-sensitive* algorithm follows the control flow order of statements within a method, and computes different solutions for a variable at distinct program points. A *context-sensitive* algorithm considers (sometimes approximately) only interprocedurally *realizable paths* [SP81, LR91, RHS95]: paths along which calls and returns are properly matched, while a *context-insensitive* algorithm does not make this distinction.

Existing algorithms for points-to analysis vary in their flow and context sensitivity; they compute approximate solutions which are supersets of the precise points-to solution. The least expensive, but most imprecise, are the flow- and context-insensitive approaches [Wei80, And94, Ste96, ZRL96, SH97]. In contrast, the flow- and context-sensitive techniques [LR92, CBC93, MLR+93, EGH94, Deu94, CHS95, WL95, Ruf95, PR96] are the most precise but also the most expensive (in time and memory). Although most of these have been used for alias or points-to analysis of C, they can be adapted for points-to analysis of C^{++}without exceptions.

The precision of the solution computed for points-to analysis directly affects its utility in applications. However, flow- and context-sensitive points-to analysis of C^{++}/Java is difficult, due to dynamic dispatch, objects containing pointers to subobjects (i.e., recursive types), exceptions, and the many invocation contexts per method. In addition, flow- and context-sensitive algorithms are memory intensive; they frequently run out of memory while analyzing even moderately sized programs. These difficulties were addressed in the design (and implementation) of *RCI*, a modular technique explained here as applied to points-to analysis for a simple object-oriented language that captures the essential features of $C^{++}$and Java (except threads). By the term *modular*, we mean a technique for whole program analysis which never requires the entire program source to be in memory at any one time. Since much object-oriented code is written as libraries, a goal of *RCI* is to analyze incomplete programs. An additional goal for *RCI* is that it maintain a sufficient degree of precision in its data-flow information for intended applications.

Intuitively, we analyze each method assuming unknown initial values for parameters and globals at method entry. The key insight is to obtain a summary function for the data-flow effect of method execution by bottom-up inference of the *relevant* conditions on the unknown initial values. These conditions capture only the relevant contexts for a method, making this approach feasible, although the summary function can be used in any context. Previous techniques [LR92, Ema93, Ghi96, WL95], have incorporated memoization of the data-flow solution associated with a particular calling context (or set of contexts).

In our approach, method bodies are analyzed first separately from calling context information. The effects of

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a method on points-to analysis information is calculated, sometimes dependent on certain conditions on the incoming unknown initial values for parameters and globals. Rather than calculate all possible conditions, the algorithm calculates only those conditions which may affect points-to information, by inferring them from the code of the method and those other methods it may invoke directly or indirectly during its lifetime.[1] Care is taken to observe those object fields actually *used* by this method directly or indirectly through calls; conditions are inferred for only those fields which are *used* in this sense. The results of these calculations are twofold: (i) a points-to solution at each node of the method and (ii) a summary transfer function for the method, a function expressing method invocation effects on the points-to solution; both are parametrized by the unknown initial values and the conditions on these values. Summary functions for callees are calculated before those of their callers. Then, points-to information is propagated into a method from its callers, with actual-parameter bindings accounted for. Recursion requires simultaneous handling of calls in the same strongly connected component (SCC) of the program call graph. Finally, the required points-to solution at a statement is computed on demand by instantiating the unknown initial values using the points-to information at a method entry.

This entire calculation is carefully staged, so that the entire program source need not be in memory at any one time, because the calculation can be done separately on the SCC's of the program call graph.

Our algorithm is the first *modular*, flow- and context-sensitive algorithm for points-to analysis of programs written in a realistic subset of C^{++}/Java. Moreover, although in this paper we explain *RCI* in the context of points-to analysis, the technique can be extended to to solve other data-flow analysis problems.

The main results in this paper are:

- An algorithm for modular points-to analysis of programs written in a simple object-oriented language that captures the essential features of $C^{++}$and Java (except threads).
- An approach to handle exceptions in the context of *modular analysis*.
- An approach to analyze a special class of incomplete programs, such as libraries.
- The definition of a new coverage metric and a discussion of how *RCI* can help in generating relevant test cases for the unit testing of libraries.
- Empirical evidence for the effectiveness of *RCI* as applied to points-to analysis of C^{++}.

The rest of this paper is organized as follows. Section 2 contains some technical definitions needed to explain *RCI*. Section 3 presents *RCI* for points-to analysis. Section 4 explains how *RCI* handles exceptions. Section 5 explains how *RCI* analyzes libraries and how the information computed by *RCI* can be used for generating relevant test cases for libraries. Section 6 contains empirical evidence for the effectiveness of *RCI*. Section 7 compares *RCI* with other approaches. Finally, Section 8 summarizes the main points of this work and discusses directions for future work.

2 Definitions

This section presents many technical definitions needed to explain the algorithm and also delimits the subsets of C^{++} and Java that are handled.

RL. For the ease of presentation, in this paper we describe *RCI* for a simple object-oriented language **RL** that has the essential features of $C^{++}$and Java (except threads). This allows us to simplify the presentation while demonstrating the interesting parts of our algorithm. *RL* is defined in Figure 1[2]. It includes single inheritance, dynamic dispatch, recursive types, exceptions and pointer assignment statements with a single level of dereferencing.[3] The semantics of exceptions is same as in Java. The semantics of other constructs is same as in C^{++}. *RL* excludes multiple inheritance, an explicit address operator (i.e., pointers to the stack), the C^{++}reference type, function pointers, data members of structure types (note this does not exclude data members of pointers to structure types), general pointer assignment statements and arrays (array elements are mapped to a single representative element). Many of these constructs can be easily accommodated by extending the algorithm; we will briefly indicate these extensions where relevant.

If the algorithm is understood fully for *RL*, then handling of most of C^{++} and Java (except threads) requires handling of details but not any changes to the fundamental ideas of the algorithm. The subset of C^{++} that can easily be handled by *RCI* excludes arbitrary casting, uninstantiated templates, pointers to data members and pointers to member methods (which are different from ordinary function pointers). Up-cast of a derived class to a base class and down-cast of a base class to a derived class can be handled. Our implementation of the algorithm is consistent with this bigger subset, only excluding in addition, multiple inheritance and exceptions.

In addition, *RCI* can essentially handle Java without threads. Under certain circumstances, we have to exclude some other features: Since, in Java, *finalizers* are invoked non-deterministically during garbage collection, we exclude *finalizers* that modify pointers. *RCI* can handle only those *finalizers* that can be ignored for points-to analysis. We exclude static initializations that depend upon the order in which files are loaded. *RCI* can handle only those static initializations for which it is safe (with respect to points-to analysis) to consider them to be executed in program source order (or any other order specified by the user) at the start of program execution. Finally, we exclude classes whose code is constructed on the fly and is not known statically.

Precision and Safety. At each program point, points-to analysis calculates the set of objects to which a pointer may point during some execution. Any static analysis technique

[1] This approach is analogous to the way pointer aliases are calculated in [LR92]. The possible reaching aliases (RA's) used in the aliasing algorithm are those which actually are propagated to procedure entry, rather than all those aliases which possibly might be propagated to the entry. Likewise, *RCI* solves for conditions that potentially affect the data-flow solution, rather than for all conditions.

[2] *Expr* is any side-effect-free expression that does not have any function call and that can be ignored for points-to analysis. *[pattern]* means at most one occurrence of the *pattern*. $\{pattern\}^+$ means one or more occurrences of the *pattern*. $\{pattern\}^*$ means zero or more occurrences of the *pattern*. $\{a \mid b\}$ means *a* or *b*. The terminal symbols are underlined.

[3] lhs = rhs, where lhs is p or p→f1 and rhs is q or q→f2.

```
Program ⇒ {Class | Proc}+
Class ⇒ class ClassName [: public (ClassName | Exception)]
  { {DataMember| Method}+ }
DataMember ⇒ Protection [static] Type FieldName ;
Type ⇒ ClassName * | PremitiveType
PremitiveType ⇒ int | char | float | bool
Method ⇒ Protection [static | virtual] [void | Type]
  MethodName ( [Param RestParam*] ) { Body }
Param ⇒ Type VarName
RestParam ⇒ , Param
Proc ⇒ {void | Type} ProcName ( [Param RestParam*] )
  { Body }
Body ⇒ Decls Stmnt+
Decls ⇒ Decl*
Decl ⇒ Type VarName ;
Stmnt ⇒ AssignmentStmnt | NewStmnt | Call | If |
  While| ReturnStmnt | Break | Continue | Try | Throw | ;
ReturnStmnt ⇒ return VarName ;
Break ⇒ break [Label] ;
Continue ⇒ continue [Label] ;
AssignmentStmnt ⇒ Lhs = Rhs ;
Lhs ⇒ VarName | VarName→FieldName
Rhs ⇒ VarName | VarName→FieldName | 0 | Expr
Call ⇒ [VarName = ]
  {VarName→MethodName|MethodName
  |ClassName::MethodName|ProcName}
  ( [VarName RestVar*] ) ;
NewStmnt ⇒ VarName = new ClassName ( [VarName RestVar*] ) ;
RestVar ⇒ , VarName
If ⇒ if ( Expr ) {Stmnt+ } [else { Stmnt+ } ]
While ⇒ [Label :] while ( Expr ) { Stmnt+ }

Throw ⇒ throw VarName ;
Try ⇒ try { Stmnt+ } Catch* [Finally]
Catch ⇒ catch ( ClassName * VarName ) { Stmnt+ }
Finally ⇒ finally { Stmnt+ }

VarName ⇒ Name
FieldName ⇒ Name
ProcName ⇒ Name
MethodName ⇒ Name
ClassName ⇒ Name
Label ⇒ Name
Protection ⇒ public : | protected : | private :
```

Figure 1: Grammar for *RL*

for points-to analysis needs to represent a potentially infinite number of heap-allocated, run-time objects with a finite number of names. Following [LR92, CBC93, MLR+93, WL95, Ste96], we represent all the run-time objects created at a program point n with the single name $object_n$, a **heap-name**. A pair $\langle ptr, object_n \rangle$ belongs to the *precise* solution of points-to analysis at a program point m if and only if there exists an execution path from the *start* node of the program to m, such that when this path is executed, at m the pointer *ptr* points to an object created at program point n. A **safe** solution of points-to analysis is one that is a superset of the *precise* solution. Our modular points-to analysis algorithm calculates a safe solution.

Data Representations. *RCI* stores the data-flow facts together with corresponding alias and type context conditions (inferred by the analysis) in data-flow elements or **dfelms**. The initial phase of the algorithm examines each method, presuming unknown initial values for each of its parameters and any global variables. Depending on the pointer assignments in that method, these unknown initial values may appear as variable names in the points-to relations and/or in corresponding conditions associated with such relations.

We use var_{init} to represent the unknown initial value for the global or parameter *var*. Note that var_{init} denotes the unknown initial object to which *var* points and not the address of that object. $var_{init}.next_{init}$ represents the unknown initial value to which $var \rightarrow next$ points. $var_{init}.next_{init}.next_{init}$ represents the unknown initial value to which $var \rightarrow next \rightarrow next$ points. Obviously, in the presence of recursive types, the number of unknown initial values accessed in a method could be unbounded. To overcome this problem, unknown initial values are mapped into a finite number of sets. All the elements in a set are represented by a single, representative name. *RCI* uses patterns in the *access paths* [Deu94] of unknown initial values to form these sets (see Appendix D for details).

For points-to analysis, each *dfelm* propagated by *RCI* has the form:

$\langle$*relevant context, points-to*$\rangle$,

where *points-to* represents a pair of the form:

$\langle$*var, object*$\rangle$.

var is a local pointer variable, a global pointer variable, an unknown initial value's field of pointer type or a heap-name's field of pointer type; and *object* is an unknown initial value or heap-name. *object* can also be *null*, which is treated as a special heap-name.

A *relevant context* has the form:

$\langle$*alias context, type context*$\rangle$.

An *alias context* is *empty* or it is a conjunction of potential aliases and potential non-aliases between unknown initial values. Each **potential alias** has the form:

$(uiv_1$ **eq** $uiv_2)$

and each **potential non-alias** has the form:

$(uiv_1$ **neq** $uiv_2)$.

Here uiv_1 and uiv_2 are unknown initial values. These potential aliases and potential non-aliases are inferred by the algorithm during analysis; some *dfelms* are dependent on specific conditions of this sort.

A *type context* is a conjunction of type constraints or *empty*. Each type constraint has the form:

(type(uv) $\in$ *A::x)*,

where *uv* is an unknown initial value, *A* is a class and *x* is a dynamically dispatched method defined in *A* (a virtual method in C^{++}). *A::x* represents the set of classes containing *A* and all the subtypes of *A* for which a virtual invocation of method *x* will be resolved to the definition of method *x* in class *A*, *A::x*. The constraint means that the associated *dfelm* is valid only in those contexts in which the concrete, run-time type of *uv* (not the declared type) belongs to *A::x*.

As with reaching aliases in our earlier work on flow-sensitive aliasing[LR92], the relevant context of a *dfelm* provides a good approximation for a calling context. This allows data-flow information to be propagated along good approximations of realizable paths.

3 Modular Points-to Analysis

In this section, we explain *RCI* as applied to points-to analysis.

3.1 Four Phases of *RCI*

RCI is an iterative worklist algorithm that is flow- and context-sensitive. *RCI* takes as input a statement-level interprocedural control flow graph or **ICFG** [LR92]. From this, first an initial approximate call graph is formed and then decomposed into strongly connected components (SCC's). The following phases are performed using the SCC condensation (**SCC-DAG**).

- **Phase 0:** In this phase, *RCI* constructs a safe overestimate of the call graph called the **initial call graph** by resolving dynamically dispatched calls using hierarchy analysis [DMM96].[4] Then *RCI* uses a linear-time algorithm[CLR92] to construct the SCC-DAG of the initial call graph. Note that the initial call graph need not be precise, it only needs to be a safe overestimate; the precision of any safe initial call graph only affects the *efficiency* of *RCI*, and not the *safety* of the computed solution. The initial call graph can be made more precise (e.g., by using [BS96]); however, in practice we have found hierarchy analysis to be adequate.

- **Phase I:** *RCI* traverses the SCC-DAG in a reverse topological order (bottom-up) and analyzes each method assuming parameters and global variables have unknown initial values. For each method *RCI* computes, in terms of unknown initial values, a safe approximation to the method's complete transfer function for pointers. We will call this approximation, the **summary transfer function**. The summary transfer function of a method *M* is the set of *dfelms* that reach the exit node of *M* and that do not represent values of local variables of *M*. This function summarizes the possible effects of method invocation on *dfelms*. The summary transfer functions of methods in the same SCC have cyclic dependences, so they are computed simultaneously by fixed point iteration. In contrast, the summary transfer functions of methods in different SCC's have hierarchical dependences (or no dependence at all), and hence are computed by bottom-up traversal on SCC-DAG, without iteration. After Phase I, the summary transfer function of each method and the points-to solution at each node in the method are expressed in terms of the *dfelms*, defined in Section 2, which may contain unknown initial values.

- **Phase II:** *RCI* traverses the SCC-DAG in a topological order (top-down) and propagates *concrete values* of unknown initial values to the *entry* nodes of methods. This phase involves only the entry nodes and call nodes, and as our empirical results show, it is extremely fast. *RCI* considers only *reachable* methods during this phase.

- **Phase III:** This phase involves only nodes that are not entry nodes. In this phase, the unknown initial values in the *dfelms* computed during Phase I are instantiated by their concrete values computed in Phase II. This phase is completely *demand-driven*, and needs to be performed only at those nodes where the final solution is needed. After this phase, the points-to solution at each node is expressed entirely in terms of program variables and heap-names.

[4] In the initial call graph, function pointer call targets can be approximated by those functions whose addresses have been stored and whose signatures match with the type of the called function.

```
for each SCC in reverse topological order
  mark-reachable()
  initialize-worklist()
  process-worklist()
```

Figure 2: Phase I

During the construction of the initial call graph using hierarchy analysis, each method needs to be in memory once. After this, each node of SCC-DAG (and hence each method) needs to be in memory only three more times, once during each of the phases I, II and III. The rest of the time, only a method's summary transfer function or the Phase II solution at the entry node of a method needs to be in memory. Hence, this is a *modular* approach and requires less memory than other whole-program-analysis techniques, in which a method cannot be moved out of memory without the possibility of it being needed again, until the final solution is computed. In these techniques, if the whole program is not kept in memory, there is no *a priori* constant bound on the number of times a method needs to be moved into or out of memory.

Of course, in the worst case, the entire initial call graph may be a single SCC and *RCI* may need to keep the whole program in memory. However, our empirical results show that SCC's are quite small in practice and *RCI* is able to analyze almost a method at a time. In some very specific domains (e.g., recursive descent parsing), SCC's may be occasionally large, however, even in these cases the entire initial call graph is unlikely to be a single SCC. For example, in parsing, the SCC's of methods dealing with statements are likely to be different from the SCC's of methods dealing with types.

3.2 Phase I

A high-level pseudocode for Phase I is given in Figure 2.

mark-reachable (Appendix G) marks reachable *ICFG* nodes in the current SCC using a simple algorithm. It assumes the entry node of each method in the current SCC is reachable.

In Phase I, *RCI* analyzes each method using the unknown initial values of parameters and global variables. By analyzing a method, it infers the relevant potential aliasing between unknown initial values and their relevant potential concrete (not declared) types. *RCI* computes a *dfelm* conditioned on such potential aliases (i.e., *dfelm*'s alias context) and concrete types (i.e., *dfelm*'s type context).

Propagation in this phase occurs in reverse topological order on the SCC-DAG; the objective is to calculate the points-to solution at each node in terms of unknown initial values and to calculate a summary transfer function for each method in the program, so it can be used to compute the summary transfer function of the callers of the method. Intuitively, a *dfelm* d at the exit node of a method M is valid at a call site that invokes M if and only if all the conjuncts of the relevant context of d are *true* at the call site. In general, at a call site that invokes M, one of the following three things happens for a conjunct of the relevant context of d: the conjunct evaluates to *true*, it evaluates to *false*, or it is translated into a similar conjunct involving the unknown initial values of the caller.

For propagation of *dfelms* between methods in different SCC's, no iteration is necessary. Using actual-parameter bindings, the summary transfer function of the called method can be used to calculate the *dfelms* returned from the call.

However, to obtain the proper propagation of *dfelms* through a non-trivial SCC, it is necessary to propagate *dfelms* on the graph of the SCC itself. Since there are cyclic dependences, iteration must be performed until a fixed point is reached. During this iteration, only a partial summary transfer function may be available at a method's exit node; this is used when processing same-SCC callers of this method. Whenever a new *dfelm* is added to the partial summary transfer function of a method, the same-SCC callers of this method are informed about this *dfelm*, so that corresponding call sites can process this new *dfelm*.

The calculation of the points-to solution in terms of unknown initial values and the calculation of the summary transfer function of a method are accomplished by the propagation of *dfelms* from method entry to exit using a worklist algorithm. Initially the solution at each node is empty; it grows monotonically as new *dfelms* are added. The code of the method is represented in the ICFG. Points-to information which reaches an ICFG node serves as input to its node transfer function which embodies the data-flow effect of the semantics of code corresponding to that node. We will specify the statement transfer functions for points-to analysis at pointer assignments and call statements. In the next section, we present examples showing these transfer functions and how relevant contexts are manipulated; we give pseudocode for Phase I in Appendix A.

3.2.1 Examples of *dfelm* Propagation

Figure 3 shows the final Phase I solution computed by *RCI* at the top of each of the statements 1, 2, 3 and 4 (i.e., program points 1,2,3,4). The *dfelms* at program point 1 say that all the variables have the same values as at the entry node of *method1*. Since statement 1 assigns the value of *a3* to the *next* field of the value of *a1*, **2:1**, the first *dfelm* at program point 2, says that the *next* field of the unknown initial value of *a1* points to the unknown initial value of *a3*. The relevant context ⟨*empty, empty*⟩ means that both alias context and type context are *empty*, and therefore this *dfelm* is valid in all contexts.

At an assignment statement, for each set s of *dfelms* that gives the values of the left and right hand sides, the relevant context of the corresponding *dfelm* (implied by s) resulting from the assignment is the conjunction of the relevant contexts of the *dfelms* of s. For statement 1, there is only one such s which consists of **1:1** and **1:3**, and the conjunction of the relevant contexts of the *dfelms* of s is ⟨*empty, empty*⟩. If the unknown initial values of *a1* and *a2* are the same, statement 1 also modifies the *next* field of the unknown initial value of *a2*. The *dfelms* **2:2** and **2:3** keep track of this potential modification. **2:2** is applicable to only those contexts in which $a1_{init}$ and $a2_{init}$ are equal, while **2:3** is applicable to only those contexts in which $a1_{init}$ and $a2_{init}$ are not equal. The *dfelms* **3:1** and **3:2** are implied respectively by **2:2** and **2:3**.

RCI does not consider $a3_{init}.next$ for potential modification at statement 1 because it is not used in *method1*. *RCI* generates fields of unknown initial values *lazily* (as explained in Section 3.2.3), only inferring *relevant* potential aliases and potential non-aliases. For example, the relationships of $a3_{init}$ with $a1_{init}$ and $a2_{init}$ are not relevant to *method1* because first, $a3_{init}.next$ is not used in *method1* and second, $a2_{init}.next$ is not directly modified in *method1*.

Statement 3 is a dynamically dispatched call site. It can invoke either *A::choose* or *C::choose* depending upon the concrete type of $a1_{init}$. As a result, *RCI* computes the values of *global2* at program point 4 conditioned on the potential concrete types of $a1_{init}$. The *dfelm* **4:1** says that *global2* points to $a1_{init}$ in those contexts in which the concrete type of $a1_{init}$ belongs to the set of types represented by *A::choose* (i.e., $\{A,B\}$). The meaning of **4:2** is similar. As before with aliases, *RCI* only infers relevant potential concrete types; for example, the concrete types of $a2_{init}$ and $a3_{init}$ are not relevant to *method1*.[5] *RCI* does not kill any *dfelm* at a call site unless the *dfelm* represents a value of the variable in which the result of the call is stored, so the *dfelms* {**1:1** to **1:3**, **2:1** to **2:3**, **3:1**, **3:2**} go across the call in statement 3.

The summary transfer function of *method1* consists of the *dfelms* at program point 4 except **1:1**, **1:2** and **1:3**, as these three *dfelms* represent values of variables local to *method1*.

Now consider the call sites that invoke *method1* from another *static* method *method2* of *test*.

```
static void method2( A* a4, A* a5, A* a6 ) {
  A *p, *q;
  5:  method1(a4,a5,a6);
  6:  p = new B(); 7:  q = new C();
  8:  method1(p,q,q);
  9:  }
```

At call site 5, *RCI* translates the *dfelms* at the exit node of *method1* by replacing the unknown initial values of *a1*, *a2* and *a3* with their values at call site 5 (i.e., the unknown initial values of *a4*, *a5* and *a6* respectively). Note: this is propagating data-flow information from callee back to caller. For example:

- ⟨⟨($a1_{init}$ **eq** $a2_{init}$),*empty*⟩, ⟨*global1*, $a3_{init}$⟩⟩
  is translated to
  ⟨⟨($a4_{init}$ **eq** $a5_{init}$),*empty*⟩, ⟨*global1*, $a6_{init}$⟩⟩.
- ⟨⟨*empty*,(*type*($a1_{init}$) ∈ *C::choose*)⟩, ⟨*global2*, $a2_{init}$⟩⟩
  is translated to
  ⟨⟨*empty*,(*type*($a4_{init}$) ∈ *C::choose*)⟩, ⟨*global2*, $a5_{init}$⟩⟩.

[5] In C^{++}, a call through an unknown initial value of a function pointer can be handled similarly by conditioning on the relationships between the unknown initial value (of the function pointer) and the functions that can be potentially invoked, that is, functions whose addresses have been stored in a function pointer and whose signatures match with the type of the call. For libraries, we assume that no function external to the library is invoked through a function pointer in the library (see Section 5).

```
class A {
  public:  A* next; public:  A() { next = 0; };
  public:  virtual A* choose( A* a1, A* a2 ) {
    return a1;
  };
};
class B: public A {
};
class C: public A {
  public:  A* choose( A* a1, A* a2 ) {
    return a2;
  };
};
class test {
  public:  static A* global1;
  public:  static A* global2;
  public:  static void method1( A* a1, A* a2, A* a3 ) {
    1:  a1 → next = a3;
    2:  global1 = a2 → next;
    3:  global2 = a1 → choose( a1, a2 );
    4:
  };
  ...
};
```

1:1 $\langle\langle \mathit{empty},\ \mathit{empty}\rangle,\ \langle \mathit{a1},\ a1_{init}\rangle\rangle$
1:2 $\langle\langle \mathit{empty},\ \mathit{empty}\rangle,\ \langle \mathit{a2},\ a2_{init}\rangle\rangle$
1:3 $\langle\langle \mathit{empty},\ \mathit{empty}\rangle,\ \langle \mathit{a3},\ a3_{init}\rangle\rangle$
1:4 $\langle\langle \mathit{empty},\ \mathit{empty}\rangle,\ \langle \mathit{global1},\ global1_{init}\rangle\rangle$
1:5 $\langle\langle \mathit{empty},\ \mathit{empty}\rangle,\ \langle \mathit{global2},\ global2_{init}\rangle\rangle$
1:6 $\langle\langle \mathit{empty},\ \mathit{empty}\rangle,\ \langle a1_{init}.\mathit{next},\ a1_{init}.\mathit{next}_{init}\rangle\rangle$
1:7 $\langle\langle \mathit{empty},\ \mathit{empty}\rangle,\ \langle a2_{init}.\mathit{next},\ a2_{init}.\mathit{next}_{init}\rangle\rangle$

2:1 $\langle\langle \mathit{empty},\ \mathit{empty}\rangle,\ \langle a1_{init}.\mathit{next},\ a3_{init}\rangle\rangle$
2:2 $\langle\langle (a1_{init}\ \mathbf{eq}\ a2_{init}),\mathit{empty}\rangle,\ \langle a2_{init}.\mathit{next},\ a3_{init}\rangle\rangle$
2:3 $\langle\langle (a1_{init}\ \mathbf{neq}\ a2_{init}),\mathit{empty}\rangle,\ \langle a2_{init}.\mathit{next},\ a2_{init}.\mathit{next}_{init}\rangle\rangle$
2: + {**1:1** to **1:5**}

3:1 $\langle\langle (a1_{init}\ \mathbf{eq}\ a2_{init}),\mathit{empty}\rangle,\ \langle \mathit{global1},\ a3_{init}\rangle\rangle$
3:2 $\langle\langle (a1_{init}\ \mathbf{neq}\ a2_{init}),\mathit{empty}\rangle,\ \langle \mathit{global1},\ a2_{init}.\mathit{next}_{init}\rangle\rangle$
3: + {**1:1** to **1:3**,**1:5**, **2:1** to **2:3** }

4:1 $\langle\langle \mathit{empty},(\mathit{type}(a1_{init}) \in \mathit{A::choose})\rangle,\ \langle \mathit{global2},\ a1_{init}\rangle\rangle$
4:2 $\langle\langle \mathit{empty},(\mathit{type}(a1_{init}) \in \mathit{C::choose})\rangle,\ \langle \mathit{global2},\ a2_{init}\rangle\rangle$
4: + {**1:1** to **1:3**,**2:1** to **2:3**, **3:1**, **3:2** }

Figure 3: Phase I Solution

At call site 8, *RCI* evaluates the relevant contexts of the *dfelms* at the exit node of *method1* using $object_6$ as $a1_{init}$ and $object_7$ as $a2_{init}$ and $a3_{init}$. For example:

- $\langle\langle (a1_{init}\ \mathbf{eq}\ a2_{init}),\mathit{empty}\rangle,\ \langle \mathit{global1},\ a3_{init}\rangle\rangle$
  is not applicable to this call site as
  ($object_6$ **eq** $object_7$) is false.

- $\langle\langle (a1_{init}\ \mathbf{neq}\ a2_{init}),\mathit{empty}\rangle,\ \langle \mathit{global1},\ a2_{init}.\mathit{next}_{init}\rangle\rangle$
  is translated to
  $\langle\langle \mathit{empty},\ \mathit{empty}\rangle,\ \langle \mathit{global1},\ \mathit{null}\rangle\rangle$
  because ($object_6$ **neq** $object_7$) is *true* and the value of $object_7.\mathit{next}$ is *null* at program point 8.

- $\langle\langle \mathit{empty},(\mathit{type}(a1_{init}) \in \mathit{A::choose})\rangle,\ \langle \mathit{global2},\ a1_{init}\rangle\rangle$
  is translated to
  $\langle\langle \mathit{empty},\ \mathit{empty}\rangle,\ \langle \mathit{global2},\ object_6\rangle\rangle$
  because B, the concrete type of $a1_{init}$, belongs to A::choose, i.e., {A,B}.

To summarize, at each call site *RCI* stores the bindings between the actuals and the unknown initial values of the methods invocable from the call site and the relevant contexts of the *dfelms* that imply these bindings. At call sites 5 and 8, these relevant contexts happen to be ⟨*empty, empty*⟩. When the actual-to-unknown-initial-value bindings are used for replacing unknown initial values with actuals in a *dfelm d1*, each resulting *dfelm d2* is associated with a new relevant context. This context is the conjunction of the contexts of the bindings used in generating *d2* and the context of *d1*, instantiated with the actuals.

Although the example in Figure 3 has recursive types, only a finite number of unknown initial values are accessed in it, so representative names are not needed.

3.2.2 Lazy Strong Update

An assignment to a local or global variable can be *safely* killed at a subsequent assignment to the same variable. In contrast, an assignment to a field of a heap-name cannot be killed (without computing additional information), because a heap-name may represent more than one run-time object. *RCI* sometimes is able to kill assignments to the fields of unknown initial values because for a particular call to a method, an unknown initial value represents the same run-time object throughout that method execution[WL95]. Appendix C presents an example that illustrates this advantage of unknown initial values over heap-names.

When a points-to set of a pointer is of cardinality one, then *may* points-to information is effectively *must* points-to information (remember we explicitly track *null*). This forms the basis on which our algorithm performs *kills* of *dfelms*. *RCI* performs kills for the fields of unknown initial values lazily. During Phase I, when a *dfelm d* representing the value of a field of an unknown initial value reaches a pointer assignment node *n* that can kill *d* according to the current points-to solution (i.e., the current solution implies that *n must* update that field of the unknown initial value), the decision for killing *d* cannot be made immediately. Until a fixed point is reached, only a partial solution is available at *n* and the *must* update could change to a *may* update according to a potentially larger fixed-point solution. In this situation, *RCI* does not propagate *d* immediately to the successors of *n*; instead, it marks *n* and *d*. After the worklist for a SCC becomes empty, *RCI* revisits the marked nodes to check if the marked *dfelms* are killed according to the current solution. If not, it unmarks the *dfelms* that should not be killed, and restarts iteration to propagate these unmarked *dfelms*. The propagation stops for a SCC when none of the remaining marked *dfelms* require repropagation (i.e., the marked nodes kill the marked *dfelms* according to the fixed-point solution). In practice, we have found this scheme to be quite effective.

This lazy propagation is efficient because the nodes that need to be revisited belong to the same SCC and their number is bounded by the size of the SCC. As shown in Section 6, the SCC's are usually small.

3.2.3 Optimizations of Phase I

There are several optimizations used in Phase I to minimize the amount of data-flow information which needs to be propagated by the algorithm, as this directly affects its cost.

Limiting relevant context. Let *d* be a *dfelm* at the exit node of a method *M*, *C* be a call site that invokes *M* and *rc* be the relevant context of *d*. If *rc* evaluates to *true* at *C*, any relevant context *t* that is contained in *rc* (i.e., the set of conjuncts of *t* is a subset of the set of conjuncts of *rc*) also evaluates to *true* at *C*. As a result, we have the following theorem:[6]

Theorem 1 *For any dfelm with relevant context r, it is safe to replace r with a relevant context s that is contained in r.*

Due to Theorem 1, instead of the complete relevant context, *RCI* can use any subset of these conjuncts without compromising *safety*, although this may cause propagation of spurious *dfelms* to call sites where only a part of the original relevant context is valid (i.e., we are using approximate context sensitivity). Many heuristics can be used for choosing the part of the complete relevant context that is stored. At present we use a simple heuristic: if the user specifies a bound of *k* on the number of conjuncts of a specific kind, we store the first *k* conjuncts of this kind associated with a *dfelm*; the rest of the conjuncts are dropped. This bound is imposed uniformly for all *dfelms*; however, *RCI* allows different bounds for different *dfelms*.

Reduction of alias context. Let *t1* and *t2* be two classes. If *t1* and *t2* satisfy any of the following: (i) *t1* is same as *t2*, (ii) *t1* is a subtype of *t2* or (iii) *t1* is a supertype of *t2*, then *t1* and *t2* are called *compatible*. Two unknown initial values are called *compatible* if and only if their declared classes are *compatible*. For the concrete values of two unknown initial values to be the same, the unknown initial values must be compatible.[7] Thus, only compatible unknown initial values can participate in a potential alias or a potential non-alias. Whenever a field of an unknown initial value is modified, the same field of another compatible unknown initial value can also be potentially modified, and for *safety*, *dfelms* with appropriate alias contexts need to be generated to record this potential modification. However, the following three optimizations enable *RCI* to avoid the generation of some of these *dfelms* without compromising *safety*.

Lazy Generation of Fields. *RCI* considers only those fields of an unknown initial value that are used in a method *M*, either directly at pointer assignment statements in *M* or indirectly, through actual-to-unknown-initial-value bindings, at pointer assignment statements in methods invoked from *M* through a series of calls. For example, in Figure 3, the *next* field of $a3_{init}$ is not considered because it is not used in *method1*. When a field *f* of an unknown initial value is found to be used for the first time in *M*, this fact is propagated to the callers of *M* in the same SCC as *M*. If through actual-to-unknown-initial-value bindings, access to *f* causes access to other fields for the first time in same-SCC callers, facts representing access to these fields are also propagated iteratively using a worklist. At an assignment statement, the fields that are considered for potential modification and generation of alias contexts are the fields that are used in the method. For example, the *next* field of $a3_{init}$ is not considered for potential modification at statement 1 of *method1*, although $a3_{init}$ has the same type as $a1_{init}$.

[6] This is similar to the use of one reaching alias instead of a set of reaching aliases in [LR92].

[7] Multiple inheritance can be easily accommodated by extending the definition of *compatible* classes so that two classes are compatible if either one of the above three conditions hold or these two classes have a common derived class.

Write-only Fields. Consider the following method which is part of class *test* defined in Figure 3.

```
void method3( A* prm1, A* prm2, A* prm3 ) {
  1:  prm1 → next = prm3;
  2:  global1 = prm2 → next;
  3:  prm3 → next = global2; }
```

This method uses the *next* fields of the unknown initial values $prm1_{init}$, $prm2_{init}$ and $prm3_{init}$. Thus, the *next* fields of $prm2_{init}$ and $prm3_{init}$ should be considered for potential modification at statement 1. However, the *next* field of $prm3_{init}$ is only used for writing and it is never read in *method3*. As a result, at statement 1, it is unnecessary to generate a *dfelm* to represent the potential modification of the *next* field of $prm3_{init}$. This modification will be automatically seen at a call site of *method3* where the values of $prm1_{init}$ and $prm3_{init}$ are the same. On the other hand, consider a call site *C* that invokes *method3* and where it is possible for the values of $prm1_{init}$ and $prm3_{init}$ to be distinct unknown initial values. Suppose *C* is contained in method *M*. Further, let *uv1* and *uv2* be two distinct unknown initial values that are respectively values of $prm1_{init}$ and $prm3_{init}$ at *C*. If the *next* field of *uv2* is read in *M*, *RCI* will generate a *dfelm* with appropriate alias context in *M* to record the potential modification to the *next* field of *uv2* due to the modification of the *next* field of *uv1* at the statement 1 of *method3*.

Initially, *RCI* considers all the fields of unknown initial values to be *write-only*. When a field is found to be read for the first time, it is made *read/write* and a candidate for potential modification; thereafter, it will be considered for the generation of alias context.

Restricting alias context to unknown initial values. At a pointer-assignment statement, *RCI* only considers the fields of unknown initial values (not heap-names) for potential modification and the generation of alias context. This is safe because for any particular call to a method, the run-time objects represented by a heap-name in the method during Phase I are different from the run-time objects represented by any of the unknown initial values. Although a heap-name appearing in a method during Phase I could be associated with an unknown initial value of the method in Phase II, for any particular call to the method, the run-time objects represented by the heap-name in the two cases are different. Consider the following methods which are part of class *test* defined in Figure 3:

```
A* method4( A *prm ) {       void method5() {
  A *p, *q;                    A *r, *s;
  1:  p = new A();             4:  r = new A();
  2:  prm → next = 0;          5:  s = method4(r);
  3:  q = p → next;            6:  r = method4(s); }
  return p; }
```

The heap-name $object_1$ is a value of the unknown initial value prm_{init} and it is also used in *method4* during Phase I. The *next* field of $object_1$ is read at statement 3 and the type of $object_1$ is the same as the type of prm_{init}. However, at statement 2, when the *next* field of prm_{init} is modified, the *next* field of $object_1$ need not be considered for potential modification. This is safe because at statement 6, where $object_1$ is the value of prm_{init}, $object_1$ represents run-time objects created at statement 1 when *method4* is called from statement 5; and for the call to *method4* at statement 6, $object_1$ appearing in the Phase I solution of *method4* represents run-time objects created at statement 1 when *method4* is called from statement 6.[8]

3.3 Phase II

For each concrete value of an unknown initial value, computed at the entry node of a method M during this phase, *RCI* visits each of the call sites in M and does the following:

- For each dynamically dispatched call site C in M, *RCI* incrementally computes the set of methods invocable from C. Suppose the receiver at C is the value of a pointer variable p. Let S be the set of *dfelms* computed during Phase I that represent the values of p at C. *RCI* evaluates the *dfelms* of S by instantiating unknown initial values with their concrete values computed at the entry node of M. Those *dfelms* whose relevant contexts evaluate to *true* yield the concrete values of p at C, which are used to determine the set of methods invocable from C. Thus, at the end of Phase II, *RCI* produces the **final call graph** which is a significant refinement of the initial call graph. At a dynamically dispatched call site in the final call graph, the only targets considered invocable are those that are invocable using the concrete values of the receiver computed at the call site during Phase II.
- At each call site C, *RCI* uses the actual-to-unknown-initial-value bindings computed during Phase I (see Section 3.2.1) to propagate concrete values to the methods invocable from C. *RCI* evaluates the relevant contexts associated with an actual-to-unknown-initial-value binding by substituting the unknown initial values in the relevant contexts with their concrete values computed at the entry node of M. A binding is used for propagation if and only if at least one of the relevant contexts associated with the binding evaluates to *true*.

For methods in the same SCC, the concrete values are propagated iteratively until a fixed point is reached, while for methods in different SCC's, the propagation is done in a top-down manner without iteration.

In order to avoid propagation of concrete values from unreachable methods, *RCI* computes an initial set of *reachable* methods, and then incrementally expands this set during Phase II. For complete programs, the initial set consists of only *main*.[9] When *RCI* visits each node of SCC-DAG in a topological order during Phase II, it considers only those methods in the current SCC that have been marked reachable. Whenever *RCI* finds an unmarked method to be invocable from a call site in a reachable method, it marks the new method also as reachable.

For example, consider Phase II for the example given in Figure 3 and *method2* (see Section 3.2.1). For simplicity, assume that *method2* is reachable. As a result, *method1* is also reachable as it is invoked from *method2*. At call site 8, *RCI* stores an actual-to-unknown-initial-value binding between $object_6$ and $a1_{init}$, and the only relevant context of this binding is $\langle empty, empty\rangle$. Since the relevant context $\langle empty, empty\rangle$ trivially evaluates to *true*, Phase II propagates $object_6$ as a concrete value of $a1_{init}$ to the entry node of *method1*. At the call site in statement 3, the value of the receiver is the value of *a1*, and the value of *a1* is given by the *dfelm* $\langle\langle empty, empty\rangle, \langle a1, a1_{init}\rangle\rangle$. Phase II substitutes $object_6$ for $a1_{init}$ in this *dfelm* to obtain $object_6$ as a concrete value of the receiver. This implies that *A::choose* is invocable from statement 3 and the final call graph has an edge from statement 3 to *A::choose*.

[8]In C++ a pointer to a global can be created and hence a global can be a value of an unknown initial value. Therefore, a global can be used in a method either directly or through an unknown initial value, and thus a global can also be part of alias context.

[9]For incomplete programs like libraries, the initial set consists of all the methods that could be directly invoked from outside the incomplete program.

3.4 Phase III

Let n be a non-entry node in a reachable method M. If the solution of points-to analysis is needed at n, each *dfelm* computed at n during Phase I is instantiated with the concrete values computed at the entry node of M during Phase II. Those instantiations for which relevant contexts evaluate to *true* yield the solution of points-to analysis at n. After instantiation with concrete values, a *dfelm* yields a *points-to* of the form $\langle var, object\rangle$, where *var* is a local pointer variable, a global pointer variable or a heap-name's field of pointer type, and *object* is a heap-name.

Consider Phase III for the example in Figure 3 and *method2* (see Section 3.2.1). Suppose, Phase III needs to be done at program point 4. For this, each *dfelm* at program point 4 is instantiated with the concrete values computed at the entry node of *method1*. For example, when $a1_{init}$ is instantiated by $object_6$ in the first *dfelm* at program point 4, the relevant context of the *dfelm* evaluates to *true* and the *dfelm* yields the *points-to* $\langle global2, object_6\rangle$.

Demand Table. The solution computed at a program point by *RCI* is sufficient to find the values of variables or fields of unknown initial values used by a method, directly or indirectly through a call. Many important applications like static resolution of dynamic dispatch [PR96] and side-effect analysis [LRZ93] need only this information. However, certain other applications (e.g., detection of dangling pointers in C^{++}) may need values of variables or fields of unknown initial values not used by a method. In order to compute the values of such variables and fields, *RCI* stores additional information in a table called *DemandTable*. *RCI* uses the *DemandTable* to compute on demand the values of variables and fields of unknown initial values not used by a method. Further details are given in Appendix E.

3.5 Complexity

In this section, we briefly discuss the complexity of the various steps of *RCI*. We will focus on the important, dominating terms and, for simplicity, ignore less important terms. The analysis is for *RL* defined in Figure 1, except exceptions, but this does not significantly impact the final results.

Phase 0. Let $\underline{T_{max}}$ be the maximum number of targets of a call site in the initial call graph, let $\underline{N_c}$ be the total number of call nodes and let $\underline{N_{proc}}$ be the total number of procedures/methods. The complexity of Phase 0 is $O(T_{max}N_c + N_{proc})$.

Phase I. The scheme for dealing with recursive types ensures that the total number of unknown initial values and hence the total number of possible *dfelms* is finite, even if no bound is imposed on the number of conjuncts in a relevant context. Since *RCI* does only a finite amount of work for each *dfelm* at a program point and at each step *RCI* considers a new *dfelm* at a program point, *RCI* always terminates. Let the total number of unknown initial values generated by *RCI* be $\underline{N_{uiv}}$, the number of user-defined pointer variables be $\underline{N_{var}}$, the maximum number of fields of a class (including inherited fields) be $\underline{F_{max}}$, the total number of classes be $\underline{N_{class}}$, the number of *ICFG* nodes be $\underline{N_{nodes}}$, the total number of heap-names be $\underline{N_h}$ and the bound on the number of conjuncts in a relevant context be $\underline{k}$.

Let $\underline{N_{rc}}$ be the number of possible relevant contexts. N_{rc} is at most $O((pa+tc)^k)$, where $\{pa = 2N_{uiv}^2, tc = N_{uiv}T_{max}\}$. Here *pa* is an upper bound on the number of possible potential aliases and potential non-aliases, and *tc* is an upper bound on the number of possible type constraints. Let $\underline{N_{pt}}$ be the number of possible *points-tos*. N_{pt} is at most

$$O(fm * sm), \text{ where } \left\{ \begin{array}{l} fm = (N_h + N_{uiv})F_{max} + N_{var} \\ sm = N_h + N_{uiv} \end{array} \right\}.$$

Here fm is an upper bound on the number of pointers that can be the first member of a *points-to* pair and *sm* is an upper bound on the number of values for the second member. Now, let $\underline{N_{dfe}}$ be the total number of possible *dfelms*. N_{dfe} is at most $O(N_{rc}N_{pt})$. Hence the total number of possible *dfelms* is polynomial in N_{uiv},N_{var},N_h, T_{max} and F_{max}, assuming k is constant.

Now consider the work done by *RCI* at a pointer assignment node. For each *dfelm* reaching such a node, $O(nl * nr)$, where $\{nl = (N_{uiv} + N_h) * N_{rc}, nr = nl\}$, is an upper bound on the work done for *dfelms* directly generated by the pointer assignment statement. Here *nl* is an upper bound on the number of elements in *lhs_rc_loc_pairs* (see Appendix A) and similarly, *nr* is an upper bound on the number of elements in *rhs_rc_loc_pairs*. $O(N_{uiv} * cs)$, where $\{cs = N_{rc}N_{uiv}(N_{uiv} + N_h)\}$, is an upper bound on the work done in generating *dfelms* due to potential aliases. Here *cs* is an upper bound on the number of *dfelms* in *new_generated_dfes* (see Appendix A) that may generate *dfelms* due to potential aliases. Finally, $O(nl * cs)$ is an upper bound on the work done in generating *dfelms* due to potential non-aliases. Here *cs* is an upper bound on the number of *dfelms* in the current solution of the pointer assignment node that may generate *dfelms* due to potential non-aliases.

Similarly, it can be shown that at any node, for each *dfelm* reaching that node, the work done by *RCI* is polynomial in $N_{uiv}, N_h, N_{var}, T_{max}$ and F_{max} (assuming k is constant and a constant bound on the number of intraprocedural successors of a node). Let $\underline{N_{work}}$ be the maximum amount of work done by *RCI* for a *dfelm* at a program point.

Since at each step *RCI* considers a new *dfelm* at a program point, the total amount of work done by *RCI* is $O(N_{dfe} * N_{nodes} * N_{work})$. Hence the total work done by *RCI* is polynomial in $N_{uiv}, N_h, N_{var}, T_{max}, F_{max}$ and N_{nodes}, assuming k is constant.

$N_h, N_{var}, T_{max}, F_{max}$ and N_{nodes} are obviously bounded by the size of the program. However, in theoretically contrived cases, *RCI* can generate an exponential number of unknown initial values (see Appendix B). Although we never encountered this in practice, it can be easily avoided by enforcing a bound t on the lengths of *access paths* of unknown initial values (analogous to *k-limiting* [JM82]). Among the unknown initial values accessible from a root unknown initial value, all the unknown initial values of the same type and having *access paths* longer than t will be represented by the same representative name. This will ensure that N_{uiv} is polynomial in $N_{proc}, N_{var}, F_{max}$ and N_{class}.

Phase II. Let $\underline{C_{max}}$ be the maximum number of call nodes in a procedure/method. Let $\underline{N_{map}}$ be the maximum number of actual-to-unknown-initial-value mappings stored at a call node. N_{map} is at most $O(np * nr)$, where $\{np = (N_h + N_{uiv})N_{uiv}, nr = N_{rc}\}$. Here *np* is the maximum number of pairs of actuals and unknown initial values and the term *nr* counts the relevant contexts associated with each actual-to-unknown-initial-value mapping. Let $\underline{N_{prop}}$ be the maximum amount of work done by *RCI* at a call node for a concrete value of an unknown initial value. N_{prop} is at most $O(N_{map} * (ce + cp))$, where $\{ce = N_h^{2k}k * N_h, cp = N_h\}$. Here *ce* is the worst case cost of evaluating a relevant context and the associated actual. This is because each unknown initial value can have at most N_h concrete values. *cp* is the worst case cost of propagating concrete values to the entry node of a target procedure. So the total amount of work done during this phase in the worst case is $O((N_{uiv}N_h)C_{max}N_{prop})$. Here $O(N_{uiv}N_h)$ is an upper bound on the number of pairs of unknown initial values and their concrete values.

Phase III. The worst case of evaluating a *dfelm* during this phase is $O(er * ep)$, where $\{er = N_h^{2k}k, ep = N_h^2\}$. Here *er* is the worst case cost of evaluating the relevant context of the *dfelm* and *ep* is the worst case cost of evaluating the *points-to* of the *dfelm*. Hence the worst case cost of this phase is $O(N_h^{2k+2}k * N_{nodes} * N_{dfe})$.

A *single-level type* [CRL98] is either a class all of whose fields are of primitive types or it is an array of a primitive type. For programs with only single-level types and without dynamic dispatch, [CR97] shows that an algorithm similar to *RCI* computes the *precise* solution of points-to analysis in $O(n^4)$ worst-case time, where n is roughly the program size. This is under the usual assumption of data-flow analysis: all realizable paths are executable. This is an improvement over the $O(n^7)$ worst-case bound achievable by applying previous techniques [LR91, RHS95] to this case. If dynamic dispatch or fields of pointer type are allowed, points-to analysis is P-space hard [CRL98, Lan92].

4 Exceptions

In this section, we extend *RCI* for points-to analysis in the presence of exceptions because, unlike C^{++}, exceptions are frequently used in Java programs. In [CRL98, CRL97], we showed how to do points-to analysis of whole programs in the presence of exceptions. The algorithm in [CRL98, CRL97] needs to keep the whole program in memory and cannot analyze incomplete programs. Here, we extend the ideas presented in [CRL98, CRL97] for *modular analysis*. The semantics of exceptions in Java is discussed in [GJS96, CRL97].

Data Representations. In the presence of exceptions, each *dfelm* computed by *RCI* during Phase I has one of the following two forms:

1. ⟨*ECFI, relevant context, points-to*⟩
2. ⟨*relevant context, exception object*⟩

A *dfelm* of the first kind represents a value of a pointer. *points-to* is as defined in Section 3.2. **ECFI** or *exception control-flow information* is one of the following:

- *excp-type*, the run-time type of an exception object,
- *label*, when a *finally* statement is entered without any uncaught exception, the number of the statement where control should go after exit from the *finally* statement,

- *iv*, when the unknown initial value *iv* is thrown as an exception object or
- *empty*, for propagation along a path from the entry node of a method to a statement in the method such that the statement is not contained in a *finally* statement and there is no uncaught exception along the path.

Intuitively, a *non-empty ECFI* summarizes the control-flow due to uncaught exceptions and *finally* statements along paths through which a *dfelm* reaches a program point, and determines the future propagation of the *dfelm* from the program point. *labels* are needed for paths that enter a *finally* statement without an exception: (1) due to falling through the associated *try* statement or one of the *catch* statements of the *try* statement, (2) due to a labelled *break* or *continue* [GJS96] in the *try* statement or one of its *catch* statements or (3) due to a *return* in the *try* statement or one of its *catch* statements.

A *relevant context* has the form:

⟨alias context, type context, *EOTC*⟩

alias context and type context are as defined in Section 3.2. **EOTC** or *exception object type context* is *empty* or it is a conjunction of type constraints of one of the following two forms:

1. $(type(iv) \leq T)$ or
2. $(type(iv) \not\leq T)$.

The first type constraint says that the associated *dfelm* holds only in those contexts where the concrete type of the unknown initial value *iv* is class T or a subtype of *T*. While the second type constraint says that the the associated *dfelm* holds only in those contexts where the concrete type of the unknown initial value *iv* is neither *T* nor a subtype of *T*.

A *dfelm* of the second kind represents an exception object. *exception object* is either a unknown initial value or a heap-name. The definition of *relevant context* is the same as in the first case. Intuitively, *dfelms* of the second kind are needed because they determine the values of the parameters of the *catch* statements. An *exception object* is assigned to the parameter of a *catch* statement that catches the *exception object*.

As far as exceptions are concerned, compared to [CRL98, CRL97], the use of a unknown initial value as *ECFI* and *EOTC* are the new ideas, which enable *modular analysis*.

Propagation of ***dfelms*****.** As before, during Phase I, *RCI* traverses SCC-DAG in a reverse topological order and analyzes each method in terms of unknown initial values. We will use the example in Figure 4 to briefly explain how *RCI* propagates the above *dfelms* during this phase. Details are similar to those given in Section 3.2. Figure 4 shows a part of the solution computed by *RCI* before each of the statements 2, 3, 4, 5 and 6. For simplicity, we ignore the *dfelms* representing the value of *p*.

The two *dfelms* at program point 2 represent the values of *e*. Statement 2 throws the object to which *e* points as exception object. As a result, **3:1** to **3:4** are propagated to statement 3, the exit node of the *try* statement. The *dfelm* **3:1** represents a value of *e*. The *ECFI ET2* means that **3:1** reaches program point 3 from the entry node of *method8* along paths that have an uncaught exception of run-time type *ET2*. The relevant context ⟨*empty,empty,empty*⟩ means that **3:1** holds in all contexts. **3:1** is created from **2:1** when $object_1$ is thrown at statement 2. Similarly, **3:2** also represents a value of *e*. It is created when p_{init} is thrown at statement 2. p_{init} is the *ECFI* because the run-time type of the exception object depends upon the type of concrete value of p_{init} and could be any subtype of the declared type of p_{init}. When a *dfelm* that has a unknown initial value as *ECFI* is propagated from the exit node of a method to a call site that invokes the method, the values of the unknown initial value at the call site are used to determine the types of the exception. **3:5** also represents a value of *e* and the *ECFI empty* means that **3:5** reaches program point 3 along paths without any uncaught exceptions. **3:3** and **3:4** represent the exception objects thrown at statement 2. Again, the relevant context ⟨*empty,empty,empty*⟩ means that **3:3** and **3:4** hold in all contexts.

Now consider the four *dfelms* at program point 4. *RCI* uses the *ECFI*'s of the *dfelms* at the exit node of a *try* statement to decide how to propagate the *dfelms*. The *ECFI* of **3:1** is *ET2* and exceptions of type *ET2* or any subtype of *ET2* are caught by the *catch* statement. So **3:1** is propagated to the entry node of the *catch* statement and **4:1** is generated from it because the uncaught exception of type *ET2* has been caught by the *catch* statement. Similarly, since the exception object p_{init} is caught by the *catch* statement if and only if the type of the concrete value of p_{init} is *ET2* or a subtype of *ET2*, **4:2** is generated from **3:2**. **3:3** represents an uncaught exception of run-time type *ET2*. So it is propagated to the entry node of the *catch* statement, where it is used to instantiate the parameter of the *catch* statement. As a result, **4:3** is generated. Similarly, **4:4** is generated from **3:4**.

Next, consider the four *dfelms* at program point 5. **3:5** reaches program point 3 along a path without any uncaught exception. So **3:5** should propagate to program point 6. However, the *finally* statement must be executed no matter how the *try* statement terminates, with an exception or without an exception. *RCI* records the fact that **3:5** has to be propagated to program point 6 after the *finally* statement is executed as *ECFI 6* and the result is **5:1**. Similarly, **4:1** also generates **5:1** and **4:2** generates **5:2**. Since, the exception object p_{init} is not caught by the *catch* statement if and only if the type of the concrete value of p_{init} is neither *ET2* nor a subtype of *ET2*, **3:2** generates **5:3**. Similarly, **3:4** generates **5:4**.

Finally, consider the four *dfelms* at program point 6. These are generated respectively from the four *dfelms* at program point 5 in the obvious way.

A call or *try* statement inside a *finally* statement can cause *ECFI*'s to stack up. At a call node inside a *finally* statement, with each relevant context *r* of an actual-to-unknown initial value binding, *RCI* also remembers the corresponding *ECFI*'s of the *dfelms* that imply the binding and have *r* as their relevant context. When the actual is substituted for the unknown initial value in a *dfelm d* propagated from the exit node of the callee to the call site, the *ECFI*'s of the actual-to-unknown initial value binding are used as *ECFI*'s of the *dfelms* resulting from the substitution. However, if *d* already has an *ECFI* that represents an uncaught exception, then, according to the semantics of Java [GJS96], the *ECFI* of *d* overrides the *ECFI*'s of the actual-to-unknown initial value binding. This is the reason why one *ECFI* is sufficient. A *try* statement nested inside a *finally* statement is treated like a call to an anonymous procedure.

Phases 0, II and III of *RCI* are same as in the absence of exceptions. Also, Theorem 1 is true for the new definition of relevant context. Moreover, under the assumptions explained in Section 3.5, *RCI* is polynomial-time even in the presence of exceptions. The hardness of points-to analysis in the presence of exceptions is discussed in [CRL98, CRL97].

We have considered only exceptions generated by *throw* statements; since run-time exceptions can be generated by almost any statement, we have ignored them. Our algorithm can handle run-time exceptions if the set of statements that can generate these exceptions is given as an input. If all statements that can potentially generate run-time exceptions are considered, we will get a safe solution; however, this may generate far more information than what is useful.

5 Analysis and Testing of Libraries

In this section, we explain how *RCI* analyzes incomplete programs like libraries and how alias context, type context and *EOTC* computed by *RCI* can be used for unit testing of libraries.

Analysis of Libraries. Let M be a method in a library L. M is called a **public interface method** of L if and only if M is in the *public interface* of L or it overrides a method in the *public interface* of L. A *public interface method* can be directly invoked from a call site outside the library. A unknown initial value at the entry node of a *public interface method* is called an **interface initial value**. During Phase II, *RCI* treats an interface initial value like a concrete value and propagates it to the entry nodes of other methods if it is the value of an unknown initial value at a call site. During Phase II or III, when *RCI* instantiates an unknown initial value in a conjunct with an interface initial value, it makes conservative, worst-case assumptions about the interface initial value for evaluating the conjunct.

The above scheme is for those situations in which the points-to solution for a library is needed without the availability of any driver (e.g., for optimizing the library before shipping it). Another possibility is not to perform Phase II and Phase III on the library until drivers are available. In this case, different drivers can share the Phase I solution of the library and the cost of doing Phase I on the library will be amortized over the drivers. This is a big advantage because Phase I is the costliest step and Phase II and Phase III are quite fast (see Section 6).

The above schemes compute a *safe* solution (for points-to analysis) with respect to a driver program if at each virtual function invocation site in the library, one of the following two conditions holds:

- none of the receivers computed by *RCI* is an interface initial value, or
- for each interface initial value that is a possible receiver, the subtypes of the declared type of the interface initial value defined in the driver program do not override the method invocable according to the declared type of the interface initial value;

and, at each call site in the library that makes a call through a function pointer, one of the following two conditions holds:

- none of the values of the function pointer computed by *RCI* is an interface initial value, or
- an interface initial value is a possible value of the function pointer, but for each function whose address is taken in the driver program, either (i) the signature of the function does not match the signature of the functions that can be invoked through the function pointer or (ii) the address of the function is also taken in the library.

There are many such situations in practice where a library is used like a component and the above scheme is adequate. Also, in Java, *final* methods and methods of classes not in the *public interface* cannot be overridden by a driver. Moreover, the above two restrictions can be easily and automatically checked to see if the solutions computed by the above schemes are *safe* for a driver.

When an interface initial value is a possible receiver at a virtual function invocation site in a library and one of the subtypes of the declared type of the interface initial value defined in a driver overrides the method invocable from the call site according to the declared type of the interface initial value[10], *then a method that is not known while analyzing the library can be invoked from a call site in the library.* In this case, after the call site, one of the possible values of each variable or field that can be modified by the *unknown overriding method* or **UOM** is a **TUobj** (typed unknown object) of the declared type of the variable or field. There is one *TUobj* for each type. *RCI* makes conservative, worst-case assumptions about *TUobj*'s while analyzing the library. At a call site C that invokes a *public interface method* from a driver, a *TUobj* of type T maps to all the heap-names whose allocation sites are reachable from C and whose types are T or a subtype of T. Note that at a call site that can invoke a *UOM*, only global variables and fields of objects accessible through the values of the globals and parameters at the call site can be modified or read by the *UOM*. Moreover, the library writer can use his/her knowledge of the semantics of the called methods to provide stylized annotations about globals and type signatures of fields that can be modified or read by *UOM*'s. The type signature of a field f is a pair of the name of the class that defines f and f*'s* own name. Among the globals and fields of unknown initial values and heap-names that can be modified by an *UOM*, only the globals mentioned in the annotations and the fields whose type signatures are mentioned in the annotations need to point to *TUobj*'s after a call site that can invoke the *UOM*. The implementation of *TUobj*'s, the investigation of their impact on precision and the use of annotations to increase precision are part of future work.

Application to Testing of Libraries. Alias context, type context and *EOTC* computed by *RCI* can be used for generating relevant test cases for libraries. These contexts provide valuable information about how a library may be used in general; such information cannot be obtained using a whole-program-analysis technique on a particular driver and the library. These contexts suggest a new coverage measure for the unit testing of libraries:

- $k_1k_2k_3$*-level relevant context cover* or $k_1k_2k_3$*-RCC*,

which can be used in addition to other standard coverage measures [Bei90, RW85].

Let p be an *execution* path from the entry node of a *public interface method* M of a library L to a statement s of L, not necessarily in M, such that p does not return to the call site in a driver program from which p starts. Let x be a variable, a field of an unknown initial value of M or a field

[10] or an analogous situation exists for a call site that calls through a function pointer.

```
class ET1 :  public Exception { };
class ET2 :  public ET1 { };
class ET3 :  public ET2 { };
void method8 ( ET1 *p ) {
  1:  ET1 *e = new ET2();
  try {
    if ( _ ) {
      if ( _ ) e = p;
      2:  throw e;
    }
    3:
  }
  catch ( ET2 *excp ) {
    4:
  }
  finally {
    5:
  }
  6:
};
```

2:1 $\langle empty, \langle empty, empty, empty\rangle, \langle e,\ object_1\rangle\rangle$
2:2 $\langle empty, \langle empty, empty, empty\rangle, \langle e,\ p_{init}\rangle\rangle$

3:1 $\langle ET2, \langle empty, empty, empty\rangle, \langle e,\ object_1\rangle\rangle$
3:2 $\langle p_{init}, \langle empty, empty, empty\rangle, \langle e,\ p_{init}\rangle\rangle$
3:3 $\langle\langle empty, empty, empty\rangle,\ object_1\rangle$
3:4 $\langle\langle empty, empty, empty\rangle,\ p_{init}\rangle$
3:5 $\langle empty, \langle empty, empty, empty\rangle, \langle e,\ object_1\rangle\rangle$

4:1 $\langle empty, \langle empty, empty, empty\rangle, \langle e,\ object_1\rangle\rangle$
4:2 $\langle empty, \langle empty, empty, (type(p_{init}) \leq ET2)\rangle, \langle e,\ p_{init}\rangle\rangle$
4:3 $\langle empty, \langle empty, empty, empty\rangle, \langle excp,\ object_1\rangle\rangle$
4:4 $\langle empty, \langle empty, empty, (type(p_{init}) \leq ET2)\rangle, \langle excp,\ p_{init}\rangle\rangle$

5:1 $\langle 6, \langle empty, empty, empty\rangle, \langle e,\ object_1\rangle\rangle$
5:2 $\langle 6, \langle empty, empty, (type(p_{init}) \leq ET2)\rangle, \langle e,\ p_{init}\rangle\rangle$
5:3 $\langle p_{init}, \langle empty, empty, (type(p_{init}) \not\leq ET2)\rangle, \langle e,\ p_{init}\rangle\rangle$
5:4 $\langle\langle empty, empty, (type(p_{init}) \not\leq ET2)\rangle,\ p_{init}\rangle$

6:1 $\langle empty, \langle empty, empty, empty\rangle, \langle e,\ object_1\rangle\rangle$
6:2 $\langle empty, \langle empty, empty, (type(p_{init}) \leq ET2)\rangle, \langle e,\ p_{init}\rangle\rangle$
6:3 $\langle p_{init}, \langle empty, empty, (type(p_{init}) \not\leq ET2)\rangle, \langle e,\ p_{init}\rangle\rangle$
6:4 $\langle\langle empty, empty, (type(p_{init}) \not\leq ET2)\rangle,\ p_{init}\rangle$

Figure 4: Exceptions

of a heap-name such that x is of *reference type* and x is read at s when p is executed. An relevant context z is called the **contextual support** of a *points-to* $\langle x,y\rangle$ with respect to p if and only if the following two conditions hold:

1. If z holds at the entry node of *M*, then when p is executed, x points to y at s.

2. Let t be any relevant context properly contained in z (i.e., the set of conjuncts of t is a proper subset of the set of conjuncts of z). Then there exists a relevant context u such that if t and u simultaneously hold at the entry node of *M*, when p is executed, x does not point to y at s.

Intuitively, z is a *minimal* relevant context such that if z holds at the entry node of *M*, then when p is executed, at s x points to y.

Let S be a set whose each element is either *empty* or it is a conjunction of potential aliases, potential non-aliases and type constraints. S is called a $k_1k_2k_3$-*RCC* of L if and only if for each *contextual support cs* of L, there exists an element e of S such that the following conditions hold:

1. Each potential alias of e is contained in *cs*.

2. Each potential non-alias of e is contained in *cs*.

3. Each type constraint of e is contained in *cs*.

4. Either each potential alias of *cs* is contained in e or e has at least k_1 potential aliases.

5. Either each potential non-alias of *cs* is contained in e or e has at least k_2 potential non-aliases.

6. Either each type constraint of *cs* is contained in e or e has at least k_3 type constraints.

Intuitively, a $k_1k_2k_3$-*RCC* says which conjunctions of potential aliases, potential non-aliases and type constraints with at most k_1 potential aliases, k_2 potential non-aliases and k_3 type constraints are relevant for points-to associations in a library and hence should be exercised during unit testing.

RCI can be easily adapted for computing a $k_1k_2k_3$-*RCC* . This is explained in Appendix F.

The following example illustrates *relevant context cover*. *method9* is part of class *test* defined in Figure 3.

```
void method9( A *a1, A *a2, A *a3 ) {
  A *l;
  1:  a1→next = a3;
  2:  l = a2→next;
};
```

Let p be the path from the entry node of *method9* to statement 2, x be $a2_{init}.next$ and y be $a3_{init}$, then z is $(a1_{init}$ **eq** $a2_{init})$. For this method, $\{(a1_{init}$ **eq** $a2_{init})$, *empty*$\}$ is a $k_1k_2k_3$-*RCC* for any k_1, k_2 and k_3. Apart from $a1_{init}$ and $a2_{init}$, potential aliasing between any other unknown initial values is not relevant for unit testing of this method, and as a result such irrelevant aliases do not appear in any *CS*. However, aliases that are irrelevant for unit testing of *method9* may be relevant for a driver program. For example, if the values of $a1_{init}$ and $a3_{init}$ are same for a driver, $a3_{init}.next$ will be modified by *method9*. However, whether this modification has any effect on the rest of the driver program depends on the semantics of the driver and cannot be tested without the driver. For unit testing of *method9*, test cases must be generated with $a1_{init}$ and $a2_{init}$ aliased. However, generating a test case with any other unknown initial values aliased (e.g, $a1_{init}$ and $a3_{init}$) is not needed. Thus, *RCI* can help in the generation of relevant test cases.

The alias contexts, type contexts and *EOTC*'s computed by *RCI* can also be used by program understanding tools. *RCI* can help by uncovering unexpected but relevant potential aliases, potential non-aliases and type constraints. Moreover, the complexity of alias contexts, type contexts and *EOTC*'s provide a measure of code complexity. Complicated alias contexts, type contexts or *EOTC*'s point out portions of code that are hard to maintain and understand.

6 Implementation

Our implementation has been built using the PROLANGS Analysis Framework (PAF) which incorporates the Edison Design Group front end for ANSI C^{++}.[11] Our initial empirical results with modular points-to analysis are encouraging; however, this is a proof-of-concept implementation and there is scope for optimization.

Table 1 contains some characteristics of the thirteen C^{++} programs we have analyzed. These are some of the benchmarks used in [PR96, BS96, CGZ95].[12] The columns *lines*, *ICFG nodes*, *methods*, *virtual calls*, *SCC's* and *Max SCC* respectively show the number of lines of code, ICFG nodes, methods, dynamically dispatched call sites, nodes in SCC-DAG and methods in the maximum-sized SCC for each program.

Table 2 contains the timings using a Sparc-20 with 352 megabytes of memory. These timings do not include time for scanning, parsing or I/O. The column *Bounds* contains pairs (i, j) which mean that the bounds on the number of potential aliases and type constraints of each relevant context were i and j respectively.

We analyzed *richards*, *deriv1*, *FeynLib*, *opProd*, *penguin*, *Bdecay* and *electron* with more than one set of bounds. The second row of (1,1) for *richards* corresponds to an analysis that only allowed either one potential alias or one type constraint per relevant context. Different bounds yielded measurable variations in those characteristics shown in Tables 2 and 3, but there was no difference in the information reported for the two applications in Table 4. We analyzed with smaller bounds those programs with large running times for higher bounds, and studied the cost-precision tradeoff with respect to the applications reported in Table 4.

Although Phase III is demand-driven, for these experiments Phase III was performed at all non-entry nodes. The Phase I solution of a library can be shared by different driver programs. This is illustrated by the Phase I timings of *Bdecay*, *electron*, *opProd* and *penguin*. For these programs, $t_1 + t_2$ means the time for the driver code is t_1 and the time for the library code is t_2, a shared cost which is only incurred once.

Although *trees* and *employ* are small benchmarks, these are important to show the two orders of magnitude timing improvement over a previous flow- and context-sensitive, whole-program-analysis technique [PR96], which took 690 and 450 seconds for these benchmarks.

In Table 3, the columns *PA's* and *TC's* show the total number of potential aliases and type constraints generated. The column *Max Size Relevant Context* shows the maximum number of conjuncts in any non-empty relevant context. In order to get an estimate of the memory saving obtained by using summary transfer functions, we normalized the size of the summary transfer function of each method with the size of the complete Phase I solution of the method. To compute the latter, we considered only those nodes which are relevant for points-to analysis: we excluded those nodes which preserve points-to information. Then we averaged these normalized sizes over all methods in a program to compute the average size of a summary transfer function for a program; results are presented as a percentage in column *Ave Size Summ Fcn.* Another possibility would have been to compare the sizes of the summary transfer functions with the number of nodes in the corresponding methods. But this is not reasonable because the cost of reanalyzing a method depends upon the size of the points-to solution of the method and the number of nodes reachable from the method during an invocation (not just the number of nodes in the method). Our comparison gives a better indication of the reduction in memory requirement achieved, as *RCI* needs to keep in memory only the summary transfer functions of the methods called from the current SCC, rather than the complete solution of those methods.

No counts are reported for potential non-aliases because at present, our implementation does not generate any potential non-aliases. By Theorem 1, this does not affect safety of the computed solution. The impact of potential non-aliases on the precision of points-to analysis is likely to be small, because a potential non-alias (e.g., $uiv_1 \mathbf{neq} uiv_2$) must evaluate to *false* to increase precision, and this will only occur if both uiv_1 and uiv_2 map to the same unknown initial value during Phase I. If they both map to the same heap-name, the potential non-alias evaluates to *true* because the heap-name can represent more than one run-time object. We have chosen to describe *RCI* as including calculation of potential non-aliases, because we believe they can be useful for other applications (e.g., analyzing libraries, testing etc).

In order to test the quality of the solution computed by *RCI*, we used the solution for two different applications: *side-effect analysis* or **MOD** [LRZ93, SRLZ98] and *virtual function resolution* [PR96]. The results for *MOD* are shown in Table 4. The column *AvConcrete MOD* shows the average number of heap-names whose fields are modified by a pointer-assignment statement according to the Phase III solution. In object-oriented programs, since a method is usually called from many different contexts, the number of heap-names modified at a statement could be large even in the precise solution. We verified by inspection that multiple calling contexts is the reason why some of the *AvConcrete MOD* numbers are high. Therefore, we used the Phase I solution to compute the average number of unknown initial values or heap-names whose fields are modified at a pointer-assignment statement. This factors out the effect of expansion of an unknown initial value into multiple concrete values because of the invocation of a method from multiple contexts. Column *AvAbstract MOD* shows this second average over all pointer-assignment statements for each benchmark.

Table 4 and Figure 5 show our results for virtual function resolution. *RCI* is probably more precise than necessary for solving virtual function resolution for C^{++}, because it is really aimed at problems where there is more gain from flow and context sensitivity; nevertheless, our results show that there are calls for which *RCI* enables much better resolution, with concomitant opportunities for aggressive optimizations such as instruction scheduling and specialization. Column *Reachable Virtual Calls* shows how many of the virtual calls in each program, are actually reachable (with the given drivers for libraries). Column *Unique Hierarchy* shows the number of reachable calls which are uniquely resolved by hierarchy analysis. Column *Differences RCI-Hierarchy* shows the number of those reachable calls for which the number of targets found by *RCI* is less than the number of targets found by hierarchy analysis. In Figure 5 we show the number of these differing calls in each program and the sizes of the differences. As noted above, changing bounds made no difference in these results. The possible loss of pre-

[11] See **http://www.prolangs.rutgers.edu/public.html**.

[12] *trees* implements trees, *deriv1* implements arithmetic expression trees, *employ* implements a class hierarchy for different kinds of employees in a company, *richards* is an operating system scheduler, *deltablue* is a symbolic constraint solver, *sampleAdv, vmatrix* and *vvector* perform matrix computations, and *FeynLib* is a library for drawing Feynman Diagrams for which *Bdecay, electron, opProd and penguin* are drivers for different kinds of elementary particles.

cision due to imposing bounds on the number of conjuncts in a relevant context was not observed in the applications we considered. For them, using lower bounds improved running time significantly; however, for other applications and benchmarks, the situation might be different.

7 Related Work

Wilson and Lam [WL95] have also used unknown initial values in their algorithm for points-to analysis of *C* programs; however, there are significant differences between *RCI* and their approach. Their algorithm needs to know the exact alias relationships between the unknown initial values before a procedure can be analyzed. Their algorithm keeps the whole program in memory and cannot analyze incomplete programs. They construct *partial transfer functions*, while *RCI* constructs summary transfer functions that approximate complete transfer functions. As their algorithm is for *C*, they handle neither dynamic dispatch nor exceptions, however, they do handle function pointers. Sometimes, if there are very few calls to a method in a complete program and very few of the possible relevant contexts occur, using summary transfer functions instead of partial transfer functions could be more costly.

Unknown initial values are similar to *non-visible variables* used in [LR92] and *invisible variables* used in [EGH94] for summarizing the values of pointers that point to out-of-scope variables. Moreover, there are aspects of *hybrid data-flow analysis* [MR90] in which unknown initial values are used to model representative external values in a local analysis.

[FF97] presents *componential set-based analysis*. It is difficult to directly compare this work with *RCI* because [FF97] was implemented for Scheme and *RCI* is for languages like $C^{++}$and Java. In [FF97] and [FF96], there is no discussion of indirect modifications through aliases. [FF96] only discusses direct assignments to variables. Although this may not be critical for a functional language like Scheme, handling of such modifications is crucial for languages like $C^{++}$and Java. Another important difference is that [FF97] simplifies the constraint systems of each of the modules separately and then combines these simplified constraint systems to compute the solution for the whole program. In the second step, the simplified constraint systems of all the modules need to be in the memory simultaneously. In contrast, *RCI* needs only one SCC and the summary transfer functions of the methods directly called from the SCC to be in memory simultaneously. *RCI* separates the iterative computation within a module from the hierarchical propagation across the modules. [FF97] performs such iterative and hierarchical propagation for the simplified systems simultaneously, which requires that all the simplified systems be in memory simultaneously. Although [FF97] is a set-based analysis, we will use the terminology of *sensitivity* (defined in Section 1) which is usually applied to data-flow analysis to compare [FF97] with *RCI* : [FF97] is *flow-insensitive* and hence does not perform *strong updates*. [FF97] keeps context by copying a function's simplified constraints, while *RCI* keeps context using summary transfer functions computed in terms of unknown initial values.

Concrete type inference and call graph construction for object-oriented languages are subsumed by our analysis. There are many non-modular whole-program-analysis approaches for these two problems. Most use constraint-based analysis [Suz81, PS91, PC94, Age95, GDDC97, DGC98], but a few [CHS95, PR96, DMM96] are data-flow-based. Both

program	lines	ICFG nodes	methods	virtual calls	SCC's	Max SCC
trees	217	280	25	3	21	3
deriv1	192	320	27	28	23	3
employ	947	463	60	4	60	1
richards	987	918	87	82	87	1
deltablue	1509	1471	112	185	112	1
sampleAdv	3130	2203	130	9	130	1
vmatrix	3607	2783	128	9	128	1
vvector	3681	3401	138	9	138	1
FeynLib	6222	7247	229	47	219	4
opProd	6289	7340	230	47	219	4
penguin	6313	7450	230	47	219	4
Bdecay	6293	7465	230	47	219	4
electron	6321	7502	230	47	219	4

Table 1: Benchmarks

program	Ph 0	Ph I	Ph II	Ph III	Bounds
trees	0.01	0.64	0.05	0.30	(∞,∞)
deriv1	0.02	5.95	0.05	0.66	(∞,∞)
deriv1		2.42	0.06	0.52	$(\infty,1)$
employ	0.01	0.76	0.05	0.31	(∞,∞)
richards	0.02	33.50	0.16	9.99	(∞,∞)
richards		17.56	0.15	8.38	(1,1)
richards		9.87	0.17	6.22	(1,1)
deltablue	0.03	4.33	0.24	1.84	(∞,∞)
sampleAdv	0.03	4.42	0.03	0.79	(∞,∞)
vmatrix	0.04	6.98	0.12	4.07	(∞,∞)
vvector	0.05	10.27	0.16	5.87	(∞,∞)
opProd	0.08	3.21+99.52	2.74	34.46	(1,1)
opProd		5.38+437.72	1.64	68.38	(2,2)
penguin	0.09	3.98+99.52	2.14	30.29	(1,1)
penguin		5.25+437.72	1.52	51.70	(2,2)
Bdecay	0.08	4.15+99.52	2.10	29.92	(1,1)
Bdecay		7.24+437.72	1.41	51.54	(2,2)
electron	0.09	5.06+99.52	2.18	33.41	(1,1)
electron		4.25+437.72	1.50	61.59	(2,2)

Table 2: Timings in Seconds

program	PA's	TC's	Max Size Relevant Context	Ave Size Summ Fcn	Bounds
trees	4	12	1	8.7%	(∞,∞)
deriv1	0	116	3	6.4%	(∞,∞)
deriv1	0	116	1	6.3%	$(\infty,1)$
employ	0	36	1	5.8%	(∞,∞)
richards	85	5	3	9.0%	(∞,∞)
richards	84	5	2	9.0%	(1,1)
richards	77	5	1	9.0%	(1,1)
deltablue	1	12	1	3.9%	(∞,∞)
sampleAdv	11	5	2	7.4%	(∞,∞)
vmatrix	10	2	2	7.4%	(∞,∞)
vvector	10	6	2	7.2%	(∞,∞)
FeynLib	115	535	1	3.5%	(1,1)
FeynLib	245	579	2	3.7%	(2,2)
opProd	115	535	1	3.5%	(1,1)
opProd	245	579	2	3.7%	(2,2)
penguin	115	535	1	3.5%	(1,1)
penguin	245	579	2	3.7%	(2,2)
Bdecay	115	535	1	3.5%	(1,1)
Bdecay	245	579	2	3.7%	(2,2)
electron	115	535	1	3.5%	(1,1)
electron	245	579	2	3.7%	(2,2)

Table 3: Performance Data

program	AvAbstract MOD	AvConcrete MOD	Reachable Virtual Calls	Unique Hierarchy	Differences RCI-Hierarchy	Bounds
trees	1	1.54	1	0	1	(∞,∞)
deriv1	1	3.67	20	10	10	(∞,∞)
employ	1	2.60	4	0	4	(∞,∞)
richards	1.100	2.48	82	81	1	(∞,∞)
deltablue	1.062	2.50	133	132	1	(∞,∞)
sampleAdv	1	5.14	2	0	2	(∞,∞)
vmatrix	1	15.14	6	5	1	(∞,∞)
vvector	1	9.40	2	0	2	(∞,∞)
opProd	1.007	2.54	5	0	5	(1,1)
penguin	1.007	1.88	6	0	6	(1,1)
Bdecay	1.007	1.98	4	0	4	(1,1)
electron	1.007	2.17	7	0	7	(1,1)

Table 4: Applications

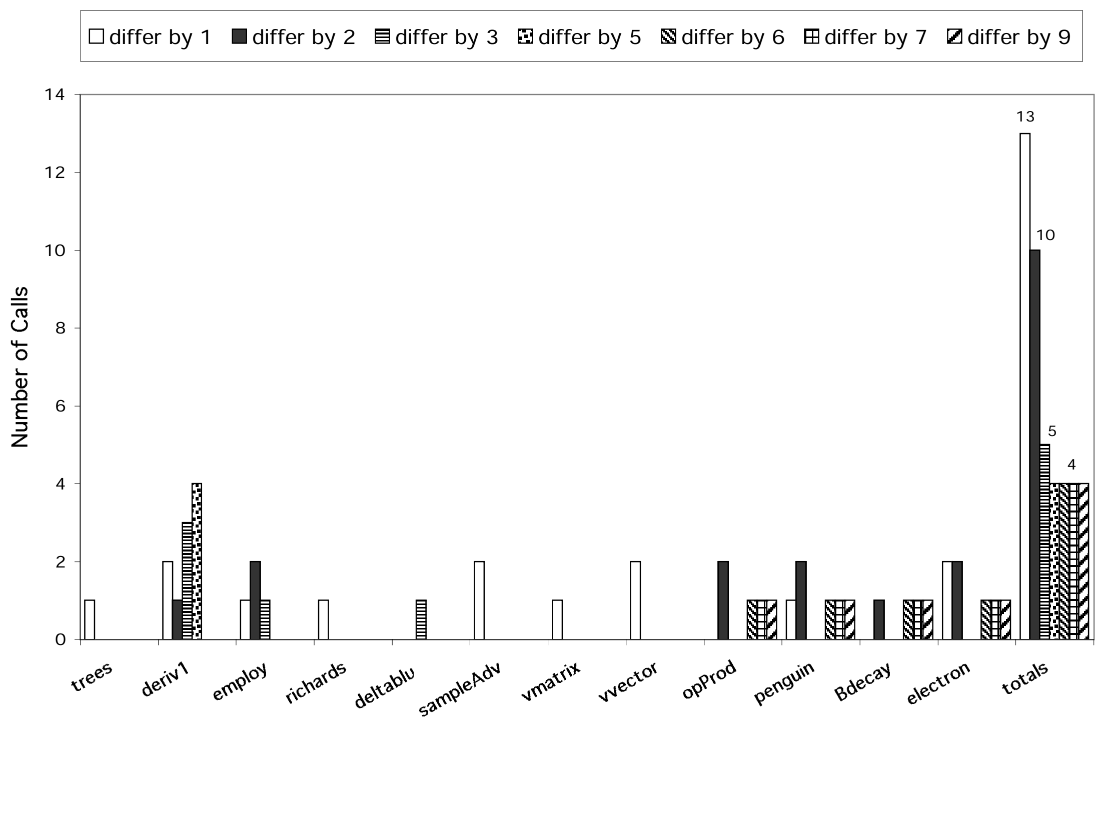

Figure 5: Differences in Precision between RCI and Hierarchy Analysis

[CHS95, PR96] are flow- and context-sensitive algorithms for C^{++}; we have compared our empirical results with those in [PR96] which were non-scalable. No implementation is reported in [CHS95]. [DMM96] presents different type analysis techniques for Modula-3, including hierarchy analysis, flow-sensitive intraprocedural type propagation and context-insensitive interprocedural type propagation. [GDDC97] presents a more complex notion of context which subsumes calling context and type context for some program variables. It presents a general framework for call graph construction of object-oriented programs. Among constraint-based approaches, [Suz81, DGC98] are flow- and context-insensitive, while [PS91, PC94, Age95] are flow-insensitive, but context-sensitive. None of the above approaches handle exceptions. [BS96] is an extension of hierarchy analysis for C^{++} and thus it is flow- and context-insensitive.

8 Conclusion and Future Work

We have presented a new technique called *relevant context inference* or *RCI* for *modular* data-flow analysis of substantial subsets of C^{++}and Java. We have explained *RCI* in the context of *modular* points-to analysis. We have shown that *RCI* can do points-to analysis even in the presence of exceptions. We have also shown that *RCI* can analyze incomplete programs like libraries. Further, we have presented a new coverage measure for unit testing of libraries and shown that the information computed by *RCI* can be used for generating relevant test cases for unit testing of libraries. Finally, we have presented empirical evidence for the effectiveness of *RCI*.

We plan to run *RCI* on larger (20000 lines) programs in future. We also plan to use *RCI* for solving other data-flow analysis problems, such as finding def-use associations. We want to explore *flow-insensitive* but *context-sensitive* variants of *RCI* as well.

RCI can also help in parallelizing data-flow analysis. Phases I and II can be done in parallel on SCC's that do not have any dependence on each other. Phase III can be done in parallel for each node at which the solution for data-flow analysis is needed. Exploration of these issues is part of future work.

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A Pointer Assignment

The data-flow transfer function, *apply*, for a pointer assignment node n, is defined in Figures 6 and 7. *n.lhs, n.rhs* and *n.sol* are respectively the left hand side expression, the right hand side expression and the current solution at n. *rdfe* is a **new** *dfelm* $\langle rc, \langle x, y\rangle\rangle$ reaching the node n (i.e., $rdfe \notin n.sol$). *lhs_rc_loc_pairs* is the set of pairs of relevant contexts and objects that are modified by node n. Note that if *n.lhs* has a dereference, *lhs_rc_loc_pairs* is initially ϕ. Otherwise, it is initially $\langle\langle empty, empty\rangle, n.lhs\rangle$. Consider the following example. Let *n.lhs* be $p{\rightarrow}f1$. A *rdfe* $(\langle rc, \langle x, y\rangle\rangle)$ can imply a new *lhs_rc_loc_pair* if and only if x is p. Thus if the on top solution for this assignment was

$$\langle\langle empty,\ empty\rangle,\ \langle p,\ l\rangle\rangle$$
$$\langle\langle (p_{init}\ \mathbf{eq}\ q_{init}), empty\rangle,\ \langle p,\ p_{init}\rangle\rangle$$
$$\langle\langle empty,\ empty\rangle,\ \langle q,\ m\rangle\rangle$$

then *lhs_rc_loc_pairs* would be

$$\left\{ \begin{array}{l} \langle\langle empty,\ empty\rangle,\ l\rangle, \\ \langle\langle (p_{init}\ \mathbf{eq}\ q_{init}), empty\rangle,\ p_{init}\rangle \end{array} \right\}$$

If *n.lhs* does not have dereference (i.e., *n.lhs* is a variable name p), then *lhs_rc_loc_pairs* contains a single element $\langle\langle empty, empty\rangle, p\rangle$ and *new_lhs_rc_loc_pairs* is always ϕ. *rhs_rc_loc_pairs* is very similar to *lhs_rc_loc_pairs* except it is for *n.rhs* (i.e., the set of pairs of relevant contexts and objects that are values of *n.rhs*).

pot_aliases computes *dfelms* that need to be generated due to potential aliases. *pot_non_aliases* computes *dfelms* that need to be generated due to potential non-aliases.

The *if* statement at program point 3 checks if *rdfe* is killed by node n. *kills* returns true if *rdfe* represents the value of a pointer variable and n directly updates this pointer variable. If n does not directly kill *rdfe*, there are two cases:

1. The assignment kills *rdfe* if certain unknown initial values are not distinct. Consider when *n.lhs* is of the form $p{\rightarrow}f1$, *rdfe* represents the value of $f1$ field of an unknown initial value s and p can point to a location that can be the concrete value of s (i.e., if p is of type $A*$, A and the type of s are *compatible*). In this case, *pot_non_aliases* is called to generate *dfelms* that condition the propagation of *rdfe* across n on potential non-aliases. These *dfelms* require that s is not equal to any of the unknown initial values that are *compatible* with s and to which p currently points. If p points to an object that cannot be same as s, *rdfe* is propagated unconditionally across n.

2. *n.lhs* cannot update the location whose value *rdfe* represents. In this case, *rdfe* is propagated across n unconditionally. Here *RCI* makes conservative assumptions about heap-names. Since, a heap-name can represent more than one run-time object, *RCI* propagates *rdfe* unconditionally across n if *rdfe* represents the value of a field of a heap-name.

```
apply( rdfe, n ) {
// rdfe = ⟨rc, ⟨x, y⟩⟩ is a new dfelm
// reaching the pointer assignment node n.

old_lhs_rc_loc_pairs = lhs_rc_loc_pairs;
old_rhs_rc_loc_pairs = rhs_rc_loc_pairs;

new_lhs_rc_loc_pairs =
new rc_loc_pairs for lhs implied by rdfe;

new_rhs_rc_loc_pairs =
new rc_loc_pairs for rhs implied by rdfe;

lhs_rc_loc_pairs =
old_lhs_rc_loc_pairs ∪ new_lhs_rc_loc_pairs;

rhs_rc_loc_pairs =
old_rhs_rc_loc_pairs ∪ new_rhs_rc_loc_pairs;

new_generated_dfes = φ;
// following loop is executed only if n.lhs
// has dereference.  Otherwise new_lhs_rc_loc_pairs
// is φ.
1:  for (each ⟨rc1, u⟩ in new_lhs_rc_loc_pairs) {
  if (u != null) {
    for (each ⟨rc2, v⟩ in rhs_rc_loc_pairs) {
      rc3 = rc1 and rc2;
      rc4 = impose user defined bounds on rc3;
      // assuming n.lhs is of the form p→f1
      new_dfe = ⟨rc4, ⟨u.f1, v⟩⟩;
      new_generated_dfes ∪ {new_dfe};
    }
  }
}

2:  for (each ⟨rc2, v⟩ in new_rhs_rc_loc_pairs) {
  for (each ⟨rc1, u⟩ in old_lhs_rc_loc_pairs) {
    if (u != null) {
      rc3 = rc1 and rc2;
      rc4 = impose user defined bounds on rc3;
      if (n.lhs has dereference)
        // assuming n.lhs is of the form p→f1
        new_dfe = ⟨rc4, ⟨u.f1, v⟩⟩;
      else
        new_dfe = ⟨rc4, ⟨u, v⟩⟩;
      new_generated_dfes ∪ {new_dfe};
    }
  }
}

new_generated_dfes = new_generated_dfes ∪
pot_aliases(new_generated_dfes, n);

new_generated_dfes = new_generated_dfes ∪
pot_non_aliases(new_lhs_rc_loc_pairs, n.sol, n);

3:  if ( !kills(n, rdfe) ) {
  if (can_update(x,n)) {
    new_generated_dfes = new_generated_dfes ∪
    pot_non_aliases(lhs_rc_loc_pairs, {rdfe}, n);
  }
  else {
    new_generated_dfes ∪ {rdfe};
  }
}

return new_generated_dfes;
}
```

Figure 6: apply1

```
kills( n, rdfe ) {
// rdfe = ⟨rc, ⟨x, y⟩⟩
  if (n.lhs does not have dereference) {
    if (x and n.lhs are the same variable)
      return true;
  }
return false;
}

can_update( x, n ) {
if (n.lhs does not have dereference)
  return false;

// else say n.lhs is p→f1 and p is of type A *
if (x is s.f1 and s is an unknown initial value whose
    type is compatible with A)
  return true;
else
  return false;
}

pot_aliases( dfes, n ) {
if (n.lhs does not have dereference)
  return φ;

// else say n.lhs is p→f1
generated_dfes_pa = φ;
for (each dfe ⟨rc1, ⟨u.f1, y⟩⟩ in dfes) {
  if (u is an unknown initial value ) {
    for (each z.f1 such that z is an unknown initial value
          compatible with u and z.f1 has been found
          to be used) {
      rc2 = rc1 and (z eq u);
      rc3 = impose user defined bounds on rc2;
      new_dfe = ⟨rc3, ⟨z.f1, y⟩⟩;
      generated_dfes_pa ∪ {new_dfe};
    }
  }
}
return generated_dfes_pa;
}

pot_non_aliases( rc_loc_pairs, dfes, n ) {
if (n.lhs does not have dereference)
  return φ;

// else say n.lhs is p→f1 and p is of type A *
generated_dfes_pna = φ;
for (each dfe ⟨rc1, ⟨u.f1, y⟩⟩ in dfes such that u is a
      unknown initial value whose type is compatible with A) {
  for (each ⟨rc2, v⟩ in rc_loc_pairs) {
    if (v is an unknown initial value compatible with u) {
      if (v != u) {
        rc3 = rc1 and rc2 and (u neq v);
        rc4 = impose user defined bounds on rc3;
        new_dfe = ⟨rc4, ⟨u.f1, y⟩⟩;
        generated_dfes_pna ∪ {new_dfe};
      }
    }
    else {
      generated_dfes_pna ∪ {⟨rc1, ⟨u.f1, y⟩⟩};
    }
  }
}
return generated_dfes_pna;
}
```

Figure 7: apply2

B Complexity of *RCI*

In the following example, at program point s, p_n points to an exponential number of unknown initial values.

```
class A_1 {
  public:  A_2 *f_1;
  public:  A_2 *f_2;
};
...
class A_i {
  public:  A_{i+1} *f_1;
  public:  A_{i+1} *f_2;
};
...
class A_{n-1} {
  public:  A_n *f_1;
  public:  A_n *f_2;
};
class A_n { };
```

```
void method( A_1 *a) {
  A_1 *p_1;
  ...
  A_i *p_i;
  ...
  A_n *p_n;
  p_1 = a;
  if (_) p_2 = p_1→f_1;
    else p_2 = p_1→f_2;
  ...
  if (_) p_i = p_{i-1}→f_1;
    else p_i = p_{i-1}→f_2;
  ...
  if (_) p_n = p_{n-1}→f_1;
    else p_n = p_{n-1}→f_2;
  s:
}
```

C Strong Update

First, Consider *method6*[13] which is part of class *test* defined in Figure 3.

```
void method6( ) {
  A *p,*q,*r;
  1:  p = new A();
  do {
    2:  r = new A();
    if ( _ ) p = r;
    q = r;
  } while( _ );
  3:  p→next = new A();
  4:  q→next = new A();
}
```

At program point 3, both p and q may point to the same heap-name $object_2$; however, they may point to different run-time objects. Thus, the assignment to the *next* field of $object_2$ at statement 3 cannot be killed by the assignment to the same field at statement 4.

Next, let us consider a slightly modified version of *method6*.

```
void method7( A *prm ) {
  A *p,*q,*r;
  1:  p = new A();
  do {
    2:  r = prm;
    if ( _ ) p = r;
    q = r;
  } while( _ );
  3:  p→next = new A();
  4:  q→next = new A();
}
```

In *method7*, at program point 3, both p and q may point to the same unknown initial value prm_{init}. Although for different calls to *method7*, prm_{init} may represent different run-time objects, for a given call to *method7*, prm_{init} represents the same run-time object throughout the method. Thus, the assignment to the *next* field of prm_{init} at statement 3 can be safely killed by the assignment to the same field at statement 4.

[13] For clarity this example is not written in *RL*.

D Recursive Types

space(uv). The unknown initial values defined with respect to the unknown initial value *uv* of a parameter or a global variable comprise the *space of initial values* associated with *uv*. *space(uv)* denotes this *space of initial values* associated with *uv*. For example:

```
class E {               class D {              class F { };
  public:  F *f1;         public:  E *f;
  public:  F *f2;       };
};

proc( D *a ) {
}
```

space(a_{init}) consists of a_{init}, $a_{init}.f_{init}$, $a_{init}.f_{init}.f1_{init}$ and $a_{init}.f_{init}.f2_{init}$.

In the presence of recursive types, *space(uv)* can be infinite as shown in the following example,

```
class H {             class G {              proc( G *p ) {
  public:               public:              };
    G *parent;            H *child;
};                    };
```

where *space(p_{init})* consists of p_{init}, $p_{init}.child_{init}$, $p_{init}.child_{init}.parent_{init}$, $p_{init}.child_{init}.parent_{init}.child_{init}$ and so on.

Representative names. Since *RCI* potentially needs to represent any arbitrary element of a *space(unknown initial value)*, in the presence of recursive types, *RCI* divides an infinite *space(unknown initial value)* into a finite number of possibly intersecting subsets. All the elements in a subset are represented by a single, representative name or *rep*. Any *dfelm* involving a *rep* represents a set of *dfelm*'s containing one *dfelm* for each instantiation of the *rep* by a member of the corresponding subset.

To intuitively describe the algorithm used by *RCI* to generate a *rep*, consider a unknown initial value $p_{init}.f1_{init}.f2_{init} \ldots f_{k_{init}}$ in *space(p_{init})*. To find the corresponding *rep*, *RCI* considers this unknown initial value as a path through the infinite tree of unknown initial values which could be possibly constructed from the type of the root unknown initial value, i.e., p_{init}. Each internal node has as its children, the unknown initial values of all the fields of the type of the internal node. A path through this tree to a node is a unknown initial value of the above form. *RCI* essentially collapses specific paths in this tree which end in the same type to one path, thus handling any recursion in the type definition. There are many ways of performing this collapse; *RCI*'s choice is just one of them.

All the unknown initial values represented by a *rep* have the same type. Subsets which contain elements of the same type are distinguished by the field selectors used in the construction of their elements. As a rule of thumb, given a unknown initial value $p_{init}.f1_{init}.f2_{init} \ldots f_{k_{init}}$, *RCI* considers the tree in which this unknown initial value is a path and starting at the root, traverses the path corresponding to the unknown initial value by selecting each field f_i in turn, for $i = 1..k$. *RCI* builds the corresponding *rep* during this tree traversal by discarding repetitive subpaths it explores. Given the selected field f_i, *RCI* looks forward on the path until it finds the last field (closest to the end of the path) with the same type as the type of f_i, say it is f_s. Then *RCI* skips $f_{i+1_{init}}, \ldots, f_{s_{init}}$ and restarts traversal from f_{s+1}. This process continues until *RCI* reaches the end of the path. Now *RCI* has the *rep* for the original unknown initial value.

A *rep* of a unknown initial value $p_{init}.f1_{init}.f2_{init} \ldots f_{k_{init}}$ has the form $v_0.v_1 \ldots v_t$, where

- v_0 is p_{init} or $[p_{init}]$ and
- for each v_i, $i = 1..k$, there exist a $l \leq k$ such that v_i is $f_{l_{init}}$ or $[f_{l_{init}}]$.

Here $[x]$ means a subpath starting at x was collapsed. Note that by construction each v_i, $i = 0..t$, has a distinct type. This ensures that the number of *rep*'s is finite.

Mapping between actuals and unknown initial values. Let *approx-space(uv)* be the approximation of *space(uv)* constructed by *RCI* using *rep*'s and unknown initial values. Here *uv* is the unknown initial value of a global or a parameter at the entry node of a method *M*. Since *approx-space(uv)* contains *rep*'s, the elements of *approx-space(uv)* can form cycles. At a call site that invokes *M*, *RCI* uses the edges (i.e., associations between fields and what they point to) between the elements of *approx-space(uv)* to determine the actuals that bind to a node of *approx-space(uv)*, which could be a unknown initial value or a *rep*. Note that as described in Section 3.2.3, *RCI* only uses the edges (or fields) that are used in *M*. Starting at a binding between an actual and the root node of *approx-space(uv)*, for each binding *b* between an actual *v* and a node *s* of *approx-space(uv)*, *RCI* recursively computes the bindings between the the nodes of *approx-space(uv)* to which the fields of *s* point and the values of the corresponding fields of *v*. Since the elements of *approx-space(uv)* form a general directed graph rather than a tree, a node of *approx-space(uv)* can map to multiple actuals.

Example. The example[14] in Figure 8 shows how *rep*'s are used to represent the summary transfer function of a method.

The *dfelm* $\langle\langle$*empty, empty*$\rangle$, $\langle$*[$param_{init}$].f2*, $x_{init}\rangle\rangle$ is part of the summary transfer function of *proc1* and *RCI* propagates it first to program point 10 and from there to program point 9. [$param_{init}$] is a *rep* and the subset represented by it binds to { $object_4$, $object_7$ } at the call site 8. Figures 9 and 10 show the bindings between the elements of *approx-space($param_{init}$)* and the actuals at call site 8. As a result, this *dfelm* expands to
{ $\langle\langle$*empty, empty*$\rangle$, $\langle object_4.f2, x_{init}\rangle\rangle$,
$\langle\langle$*empty, empty*$\rangle$, $\langle object_7.f2, x_{init}\rangle\rangle$ } at program point 9.

E Demand Driven Computation

As stated in Section 3.2.3, *RCI* generates fields of unknown initial values lazily. Similarly, it also considers globals and parameters lazily, only when they are found to be used. During Phase III, at a call node *C*, for each points-to *pt* that represents a value of the field of a heap-name at *C*, the tuple $\langle pt, C\rangle$ is stored in a global table called *DemandTable*. Similarly, for each points-to *pt* that represents a value of a global variable that is not used in at least one of the methods invocable from *C*, the tuple $\langle pt, C\rangle$ is stored in *DemandTable*. The following example[15] shows how *RCI* uses *DemandTable* to compute on demand the values of variables and fields of unknown initial values not used by a method.

```
class F { };
class E {
  public:  F *field1;
  public:  F *field2;
};
```

[14] For clarity this example is not written in *RL*.
[15] For clarity this example is not written in *RL*.

```
class A1 {
  public:  B1 *f1;
  public:  B1 *f2;
  public:  A1( ) {
    f1 = f2 = 0;
  };
};
```

```
class B1 {
  public:  C1 *f1;
  public:  A1 *f2;
  public:  B1() {
    f1 = f2 = 0;
  };
};
```

```
class C1 { };
```

```
class test1 {
  public:  static B1 *x;
  public:  static void proc0( ) {
    A1 *p;
    B1 *q;

    1:  p = new A1();

    2:  q = new B1();

    3:  p→f1 = q;

    4:  p→f1→f2 = new A1();
    5:  q = new B1();
    6:  p→f1→f2→f1 = q;
    7:  p→f1→f2→f1→f2 = new A1();
    8:  proc1( p );
    9:
  }

}
```

```
public:  static void proc1( A1 *param ) {
  A1 *tmp;

  tmp = param;

  while( tmp != 0 ) {
    tmp→f2 = x;

    if ( tmp→f1 != 0 )
      tmp = tmp→f1→f2;
    else
      break;
  }



  10:
}
```

Figure 8: Collapsing

```
class test1 {
  public:  static void foo0( E *e1 ) {
    F *p;

    5:  p = e1→field1;
    6:
  };
  public:  static void foo1( E *e2 ) {
    7:  foo0( e2 );
  };

  public:  static void main( void ) {
    E *local;

    1:  local = new E();

    2:  local→field1 = new F();

    3:  local→field2 = new F();

    4:  foo1( local );
  };
};
```

Phase I computes $\langle\langle empty,\ empty\rangle,\ \langle p,\ e1_{init}.field1_{init}\rangle\rangle$ and $\langle\langle empty,\ empty\rangle,\ \langle e2_{init}.field1,\ e2_{init}.field1_{init}\rangle\rangle$ at program points 6 and 7 respectively; and $e1_{init}.field2_{init}$ and $e2_{init}.field2_{init}$ do not appear in the Phase I solutions of *foo0* and *foo1* respectively. As a result, during Phase II, the value of $e2_{init}.field2_{init}$ at the call site 4, which is $object_3$, is not propagated to the entry node of *foo1*. Instead, the tuple $\langle\langle object_1.field2,\ object_3\rangle,\ 4\rangle$ is stored in *DemandTable*. If the value of $e1_{init}.field2_{init}$ is needed at program point *5*, then, since $e1_{init}$ is $object_1$, *RCI* looks up *DemandTable* for the values of $object_1.field2$. The result of this lookup says that $object_1.field2$ points to $object_3$ at the call site 4. Now *RCI* checks if the entry node of *foo0* is reachable from call site 4. Since it is reachable, $object_3$ is a valid value for $object_1.field2$ at program point *5*.

Context-sensitive Transitive Closure Checking reachability of an entry node from a call site is easy in the absence of dynamically dispatched calls as it reduces to computing the transitive closure of the call graph. However, in the presence of dynamically dispatched calls, the transitive closure of the call graph gives an approximate, but safe solution to the above question. This is because in the presence of dynamic dispatch, reachability is context-sensitive and hence non-transitive. In order to improve precision, *RCI* uses the following scheme for computing the context-sensitive transitive closure of the final call graph. It traverses the SCC-DAG in a reverse topological order (bottom-up) and for each *reachable* method (see Section 3.3) m, it computes the set of methods (say *reach-set*) reachable from m and relevant contexts, in terms of the unknown initial values of m, under which these methods are reachable. *RCI* uses the final call graph for determining the set of methods invocable from a call site. When the set of methods reachable from a call node C in a *reachable* method m is needed, it is computed on demand using the *reach-sets* of the methods invocable from C. The relevant contexts of the elements of these *reach-sets* are translated by instantiating the unknown initial values of the methods invocable from C with their values at C. The resulting relevant contexts are then evaluated using the concrete value's computed at the entry node of m. The set of methods reachable from C is determined by those elements whose relevant contexts evaluate to *true*. The pseudocode for context-sensitive transitive closure is given in [CRL97].

Consider the example given in Figure 3. For the sake of illustration, let us assume that the final call graph has edges from call site 3 to *A::choose* and *C::choose*. Under this assumption, the *reach-set* of *method1*, say *rs*, is

$$\{\ \langle\langle empty,(type(a1_{init}) \in A::choose)\rangle,\ A::choose\rangle,\ \langle\langle empty,(type(a1_{init}) \in C::choose)\rangle,\ C::choose\rangle\ \}$$

Now, suppose the set of methods reachable from call site 5, in *method2* (defined on page 4), is needed. The only method invocable from call site 5, in the final call graph, is *method1*. Thus, the unknown initial values of *method1* in the relevant contexts of the elements of *rs* are instantiated by their values at call site 5, i.e., $a1_{init}$ is replaced by $a4_{init}$. This yields the set *rs1*:

{ ⟨⟨*empty,(type*($a4_{init}$) ∈ *A::choose*)⟩, *A::choose*⟩,
⟨⟨*empty,(type*($a4_{init}$) ∈ *C::choose*)⟩, *C::choose*⟩ }

Now $a4_{init}$ is instantiated by its concrete value's at the entry node of *method2*, and the methods reachable from call site 5 are *method1* and the methods of the elements of *rs1* whose relevant contexts evaluate to *true* after this instantiation. Next, suppose the methods reachable from call site 8 are needed. Now, $a1_{init}$ in the relevant contexts of the elements of *rs* is instantiated by $object_6$. This means that the relevant context of the first element of *rs* evaluates to *true* and the relevant context of the second element evaluates to *false*. Hence, the set of methods reachable from call site 8 is

{ *method1*, *A::choose* }.

F Testing of Libraries

If *RCI* is used for computing only points-to's, when a *dfelm* d_1 reaches a program point n, if the points-to contained in d_1 is contained in another *dfelm* d_2 that is already present in the current solution of n and the relevant context of d_2 is contained in the relevant context of d_1, d_1 is ignored. However, d_2 could have been due to a non-executable path. This may cause *RCI* not to compute *relevant context covers*. This can be avoided by propagating newly reaching *dfelm*'s like d_1 as long as their relevant contexts obey the bounds imposed on the numbers of potential aliases, potential non-aliases and type constraints. With this modification, the following assertion about *RCI* is true. We need some notations to express this assertion. Let c be any conjunction of potential aliases, potential non-aliases and type constraints and *conc(c)* be the simplified conjunction of potential aliases, potential non-aliases and type constraints obtained from c by instantiating unknown initial values by their concrete values computed during Phase II and dropping conjuncts that are true after the instantiation. If any conjunct of c is false after the instantiation, *conc(c)* is false. *conc(c)* is called *relevant* if either it involves only interface initial value's of a particular *public interface method* or it is *empty*. Let S be a set whose each element is a conjunction of potential aliases, potential non-aliases and type constraints. Let *rconc(S)* be the set whose each element is a *relevant conc(c)* of some element c of S. Let R be the set of relevant contexts of *dfelm*'s computed by *RCI*. Then,

- if *RCI* is run on a library L with the bounds on the numbers of potential aliases, potential non-aliases and type constraints in any relevant context being at least k_1, k_2 and k_3 simultaneously, then *rconc(R)* is a $k_1k_2k_3$*-RCC* of L.

G Mark-reachable

Figures 11, 12, 13, 14, 15, 16, 17, 18 and 19 define *mark-reachable*. The data-flow elements computed by *mark-reachable* have one of the following three forms:

1. ⟨*empty,reachable*⟩
2. ⟨*excp-type,reachable*⟩
3. ⟨*label,reachable*⟩

An *exception block* is the body of a try, catch, finally or method. Given a statement S, *innermost-enclosing-exception-block(S)* is the innermost *exception block* containing S.

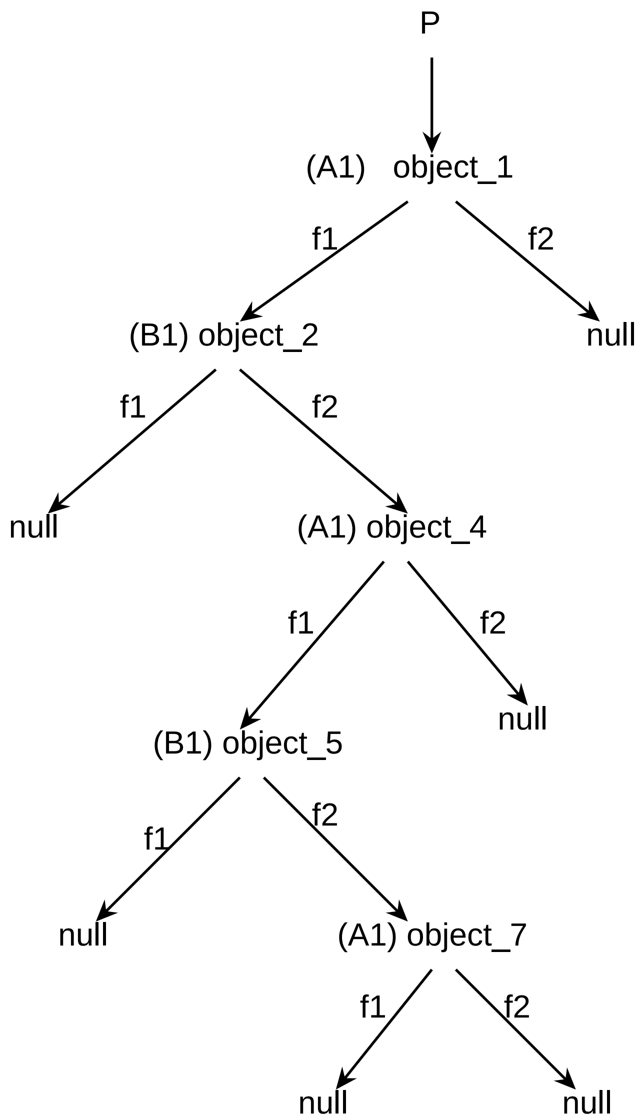

Figure 9: Actuals

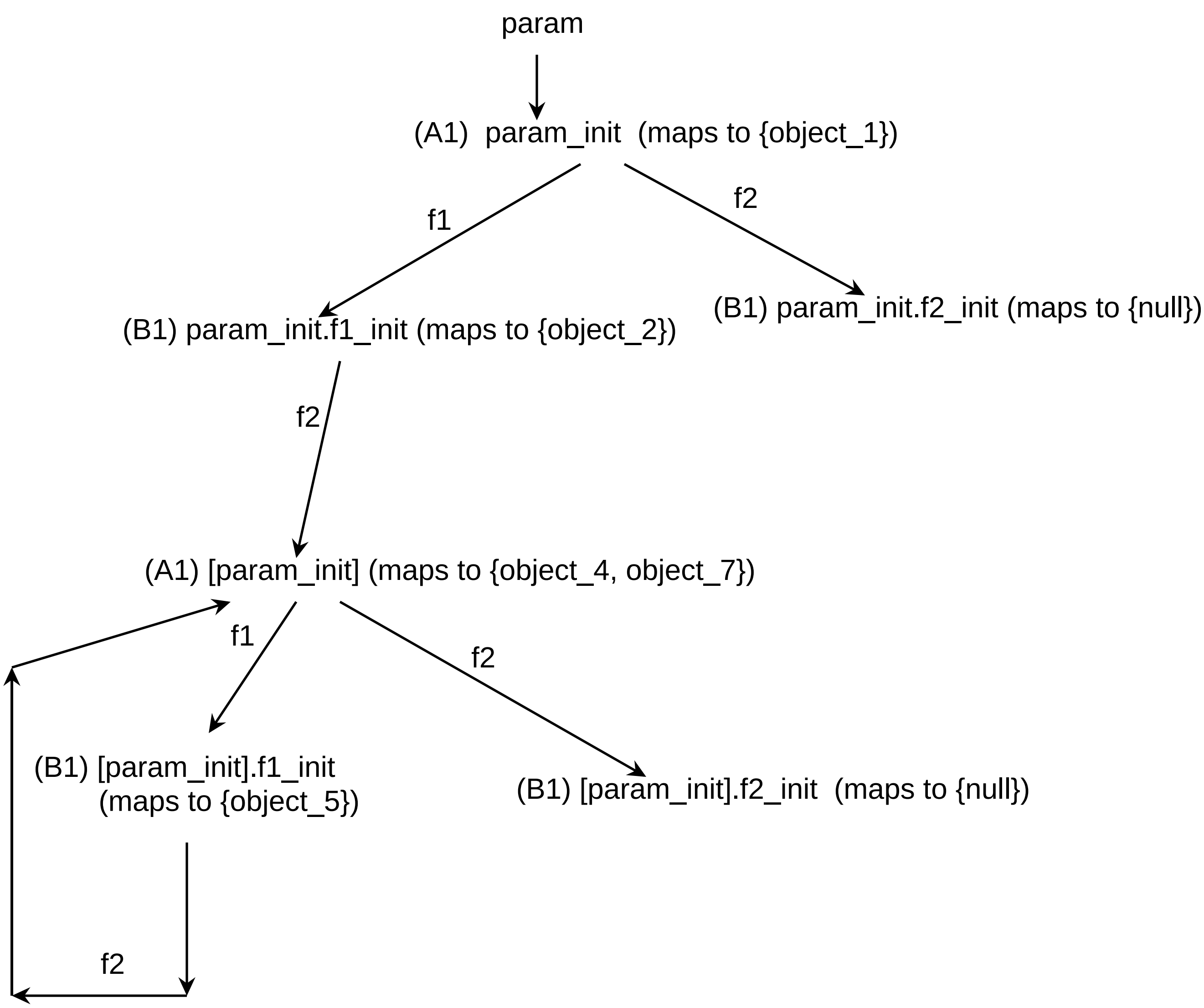

Figure 10: *approx-space*($param_{init}$)

mark-reachable() {
```
// initialize worklist
worklist = empty
for each method m in the current SCC
  wl-node = new worklist node(m.entry, ⟨empty,reachable⟩)
  add wl-node to worklist

process-worklist-mark-reachable()
for each return-site n
  n.reaching-reachability-dfes = n.reaching-reachability-dfes ∪ n.successor.reaching-reachability-dfes
}
```

Figure 11: mark-reachable

```
process-worklist-mark-reachable() {
while worklist is not empty {
  wl-node = delete node from worklist
  node = wl-node.node
  dfe = wl-node.dfe

// condition node represents the test expression of a
// if or while and new node is an object creation site
  if (node is an assignment node or condition node or new node)
    for each successor succ of node
      add-to-soln-and-worklist-if-needed-mr(succ,{dfe})

  if (node is a throw node)
    let excp-type be the declared type of the exception thrown by node
    x = exit of innermost-enclosing-exception-block(node)
    add-to-soln-and-worklist-if-needed-mr(x,{⟨excp-type,reachable⟩})

  if (node is the entry node of a try)
    process-try-entry-mr( node, dfe )
  if (node is the exit node of a try block)
    process-try-exit-mr( node, dfe )
  if (node is the entry node of a catch)
    add-to-soln-and-worklist-if-needed-mr(node.successor,{⟨empty,reachable⟩})
  if (node is the exit node of a catch)
    process-catch-exit-mr( node, dfe )
  if (node is the entry node of a finally)
    add-to-soln-and-worklist-if-needed-mr(node.successor,{dfe})
  if (node is the exit node of a finally)
    process-finally-exit-mr( node, dfe )
  if (node is the exit node of a method)
    process-method-exit-mr( node, dfe )
  if (node represents a break statement)
    process-break-mr( node, dfe )
  if (node represents a continue statement)
    process-continue-mr( node, dfe )
  if (node is a call node)
    process-call-mr( node, dfe )
  if (node is the return-site of a call node)
    // i.e., node is the successor of a call node
    process-return-site-mr( node, dfe )

  if (node represents a return statement)
    process-return-statement-mr( node, dfe )
} }
```

Figure 12: mark-reachable 1

```
process-try-entry-mr( node, dfe ) {
/*** node.try is the try statement whose entry node is node ***/
if ( innermost-enclosing-exception-block(node.try) is a finally )
  // try nested inside a finally
  // treat like a call to an anonymous procedure
  process-call-mr( node.call-node, dfe )
else
  add-to-soln-and-worklist-if-needed-mr(node.successor,{dfe})
}
```

Figure 13: mark-reachable 2

```
process-try-exit-mr( node, dfe ) {
/*** dfe.ecfi is the first element of dfe ***/
if ( dfe.ecfi represents an exception ) {
  for each catch ct associated with node
    // ct catches any exception whose type is ct.catch-type or a subtype of ct.catch-type
    if (ct.catch-type is compatible with dfe.ecfi)
      add-to-soln-and-worklist-if-needed-mr(ct.entry,{dfe})
  if (there does not exist a catch ct associated
      with node such that ct.catch-type is same as dfe.ecfi
      or ct.catch-type is a super-type of dfe.ecfi)
    // it is possible for the exception to escape
    // all catches associated with node
    propagate-to-finally-if-needed-mr( node, dfe )
  // else the exception is caught by at least one catch clause
}
else {
  propagate-to-finally-if-needed-mr( node, dfe )
} }
```

Figure 14: mark-reachable 3

```
propagate-to-finally-if-needed-mr(node, dfe) {
  if ( dfe.ecfi is empty )
    succ = successor of the try-catch-finally construct associated with node
    dfe = ⟨succ,reachable⟩
  if ( there is a finally associated with node )
    add-to-soln-and-worklist-if-needed-mr(node.finally.entry,{dfe})
  else
    if ( innermost-enclosing-exception-block(node.try) is a finally )
      process-method-exit-mr( node.method-exit, dfe )
    else
      if ( dfe.ecfi is a label contained in innermost-enclosing-exception-block(node.try) )
        add-to-soln-and-worklist-if-needed-mr(dfe.ecfi,{⟨empty,reachable⟩})
      else
        x = exit of innermost-enclosing-exception-block(node.try)
        add-to-soln-and-worklist-if-needed-mr(x,{dfe})
}
```

Figure 15: mark-reachable 3a

```
process-catch-exit-mr( node, dfe ){
  propagate-to-finally-if-needed-mr( node, dfe )
}
```

Figure 16: mark-reachable 4

```
process-finally-exit-mr( node, dfe ) {
  /*** node.try is the try statement associated with node ***/
  if ( innermost-enclosing-exception-block(node.try) is a finally )
    process-method-exit-mr( node.method-exit, dfe )
  else
    if ( dfe.ecfi is a label contained in innermost-enclosing-exception-block(node.try) )
      add-to-soln-and-worklist-if-needed-mr(dfe.ecfi,{⟨empty,reachable⟩})
    else
      x = exit of innermost-enclosing-exception-block(node.try)
      add-to-soln-and-worklist-if-needed-mr(x,{dfe})
}
```

Figure 17: mark-reachable 5

```
process-call-mr( node, dfe ) {
/*** node.ecfis contains the ECFI's of the reachability data-flow elements reaching node ***/
if (dfe.ecfi ∉ node.ecfis)
  node.ecfis = node.ecfis ∪ {dfe.ecfi}

for each method m invocable from node using hierarchy analysis
  //m.exit.reaching-reachability-dfes contains data-flow-elements that have reached the exit node of m
  if ( ⟨empty,reachable⟩ ∈ m.exit.reaching-reachability-dfes )
    add-to-soln-and-worklist-if-needed-mr(node.successor.successor,{dfe})
  for each ⟨label,reachable⟩ ∈ m.exit.reaching-reachability-dfes such that
        label is contained in innermost-enclosing-exception-block(node)
    add-to-soln-and-worklist-if-needed-mr(label,{dfe})
  if (dfe is the first data-flow-element reaching node)
    // i.e., node is found reachable for the first time
    for each dfe1 in m.exit.reaching-reachability-dfes
      if ( dfe1.ecfi is not empty and dfe1.ecfi is not a label contained
            in innermost-enclosing-exception-block(node) )
        add-to-soln-and-worklist-if-needed-mr(node.successor,{dfe1})
}
```

```
process-return-site-mr( node, dfe ) {
let dfe be ⟨ecfi,reachable⟩
if ( ecfi is empty )
  successor = ordinary successor of node
else
  successor = exit of innermost-enclosing-exception-block(node)

add-to-soln-and-worklist-if-needed-mr(successor,{dfe})
}
```

```
process-method-exit-mr( node, dfe ){
for each call site c in the current SCC that can invoke node.method {
  if (c has been found to be reachable)
    if (dfe.ecfi is empty)
      for each ecfi1 ∈ c.ecfis
        add-to-soln-and-worklist-if-needed-mr(c.successor.successor,{⟨ecfi1,reachable⟩})
    if (dfe.ecfi is a label contained in innermost-enclosing-exception-block(c) )
      for each ecfi1 ∈ c.ecfis
        add-to-soln-and-worklist-if-needed-mr(dfe.ecfi,{⟨ecfi1,reachable⟩})
    if (dfe.ecfi is an excp-type or dfe.ecfi is a label not contained in innermost-enclosing-exception-block(c) )
      add-to-soln-and-worklist-if-needed-mr(c.successor,{dfe})
}}
```

Figure 18: mark-reachable 6

```
process-break-mr( node, dfe ) {
let y = target of break
if (y is contained in innermost-enclosing-exception-block(node) )
  add-to-soln-and-worklist-if-needed-mr(y,{dfe})
else
  x = exit of innermost-enclosing-exception-block(node)
  add-to-soln-and-worklist-if-needed-mr(x,{⟨y,reachable⟩})
}
```

```
process-continue-mr( node, dfe ) {
let y = target of continue
if (y is contained in innermost-enclosing-exception-block(node) )
  add-to-soln-and-worklist-if-needed-mr(y,{dfe})
else
  x = exit of innermost-enclosing-exception-block(node)
  add-to-soln-and-worklist-if-needed-mr(x,{⟨y,reachable⟩})
}
```

```
process-return-statement-mr( node, dfe ) {
  let method-exit be the exit node of the method containing node
  if (innermost-enclosing-exception-block(node) is a method body)
    add-to-soln-and-worklist-if-needed-mr(method-exit, {dfe})
  else
    x = exit of innermost-enclosing-exception-block(node)
    add-to-soln-and-worklist-if-needed-mr(x, {⟨method-exit,reachable⟩})
}
```

```
add-to-soln-and-worklist-if-needed-mr( node, dfes ) {
for each dfe ∈ dfes
  if dfe ∉ node.reaching-reachability-dfes
    node.reaching-reachability-dfes = node.reaching-reachability-dfes ∪ {dfe}
    wl-node = new worklist node (node, dfe)
    add wl-node to worklist
}
```

Figure 19: mark-reachable 7